

# Bounding and decomposing thin analytic partial orderings \*

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## Abstract

We modify arguments in [2] to reprove extensions of two key results there in the context of bounding and decomposing of analytic subsets of Borel partial quasi-orderings.

## 1 Introduction

The following theorem is the main content of this note.

**Theorem 1.1.** *Let  $\preceq$  be a  $\Delta_1^1$  PQO on  $\omega^\omega$ ,  $\approx$  be the associated equivalence relation, and  $X^* \subseteq \omega^\omega$  be a  $\Sigma_1^1$  set such that  $\preceq \upharpoonright X^*$  is thin<sup>1</sup>. Then*

- (i) *there is an ordinal  $\alpha < \omega_1^{\text{CK}}$  and a  $\Delta_1^1$  LR order preserving map  $F : \langle \omega^\omega ; \preceq \rangle \rightarrow \langle 2^\alpha ; \leq_{1\text{ex}} \rangle$  satisfying the following additional requirement: if  $x, y \in X^*$  then  $x \not\approx y \implies F(x) \neq F(y)$  ;*
- (ii)  *$X^*$  is covered by the countable union of all  $\Delta_1^1$   $\preceq$ -chains  $C \subseteq \omega^\omega$ .*

The theorem is essentially proved in [2, 3.1 and 5.1]. Literally, only the case of  $\Delta_1^1$  subsets  $X^*$  is considered in [2], but the case of  $\Sigma_1^1$  sets  $X^*$  can be obtained by a rather transparent rearrangement of the arguments in [2]. See also [3] in matters of the additional requirement in claim (i) of the theorem, which also is presented in [2] implicitly. Our proofs will largely follow the arguments in [2], but by necessity we modify those here and there in order to streamline some key arguments. On the other hand, we substitute reflection arguments in [2] with more transparent constructions.

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<sup>1</sup> Meaning that there is no perfect set of pairwise  $\preceq$ -incomparable elements.

## 2 Notation and an important lemma

### Non-strict relations

PQO, *partial quasi-order*: reflexive ( $x \leq x$ ) and transitive in the domain;

LQO, *linear quasi-order*: PQO and  $x \leq y \vee y \leq x$  in the domain;

LO, *linear order*: LQO and  $x \leq y \wedge y \leq x \implies x = y$ ;

*associated equivalence relation*:  $x \approx y$  iff  $x \leq y \wedge y \leq x$ .

### Strict relations

*strict* PQO: irreflexive ( $x \not\leq x$ ) and transitive;

*strict* LQO: strict PQO and  $x < y \implies \forall z (z < y \vee x < z)$ ;

*strict* LO: strict PQO and the trichotomy  $\forall x, y (x < y \vee y < x \vee x = y)$ .

By default we consider only *non-strict* orderings. All cases of consideration of *strict* PQOs will be explicitly specified.

Any non-strict PQO  $\leq$  defines an associated strict one so that  $x < y$  iff  $x \leq y \wedge y \not\leq x$ . In the opposite direction, given a strict PQO  $<$ , we define an equivalence relation  $\approx$  iff  $x < z \iff y < z$  and  $z < x \iff z < y$  for all  $z$  in the domain, and then define  $x \leq y$  iff  $x < y$  or  $x \approx y$ .

### Order preserving maps

*LR (left-right) order preserving map*: any map  $f : \langle X; \leq \rangle \rightarrow \langle X'; \leq' \rangle$  such that we have  $x \leq y \implies f(x) \leq' f(y)$  for all  $x, y \in \text{dom } f$ ;

*RL (right-left) order preserving map*: any map  $f : \langle X; \leq \rangle \rightarrow \langle X'; \leq' \rangle$  such that we have  $x \leq y \iff f(x) \leq' f(y)$  for all  $x, y \in \text{dom } f$ ;

*2-ways order preserving map*: any map  $f : \langle X; \leq \rangle \rightarrow \langle X'; \leq' \rangle$  such that we have  $x \leq y \iff f(x) \leq' f(y)$  for all  $x, y \in \text{dom } f$ .

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*sub-order*: restriction of the given PQO to a subset of its domain.

$<_{\text{lex}}, \leq_{\text{lex}}$ : the lexicographical LOs on sets of the form  $2^\alpha$ ,  $\alpha \in \text{Ord}$ , resp. strict and non-strict;

*a  $\leq$ -chain in a PQO*: any set of 2wise  $\leq$ -comparable elements, i. e., LQO;

a  $\leq$ -thin set in a PQO: any set in the domain of  $\leq$  containing no perfect subsets of 2wise  $\leq$ -incomparable elements;

$[x]_{\mathbf{E}} = \{y \in \text{dom } \mathbf{E} : x \mathbf{E} y\}$  (the  $\mathbf{E}$ -class of  $x$ ) and  $[X]_{\mathbf{E}} = \bigcup_{x \in X} [x]_{\mathbf{E}}$  — whenever  $\mathbf{E}$  is an equivalence relation and  $x \in \text{dom } \mathbf{E}$ ,  $X \subseteq \text{dom } \mathbf{E}$ .

**Lemma 2.1** (Kreisel selection). *Let  $D$  be the set of all  $\Delta_1^1$  points in  $\omega^\omega$ . If  $P \subseteq \omega^\omega \times D$  is a  $\Pi_1^1$  set, and  $X \subseteq \text{dom } P$  is  $\Sigma_1^1$  then there is a  $\Delta_1^1$  set  $Y \subseteq \text{dom } P$  and a  $\Delta_1^1$  function  $F : Y \rightarrow D$  such that  $X \subseteq Y$  and  $F \subseteq P$ .*

**Proof.** The set  $X_0 = \text{dom } P$  is  $\Pi_1^1$  since  $\Pi_1^1$  is closed under  $\exists y \in \Delta_1^1$ . Therefore by Separation there is a  $\Delta_1^1$  set  $Y$ ,  $X \subseteq Y \subseteq X_0$ . By  $\Pi_1^1$  Uniformization, there is a  $\Pi_1^1$  set  $F \subseteq P$  such that  $\text{dom } F = Y$  and  $Y$  is a function. To show that  $F$  is in fact  $\Delta_1^1$  note that  $F(x) = y$  iff  $x \in Y$  and  $\forall y' \in D (y \neq y' \implies \langle x, y' \rangle \notin F)$ , which leads to a  $\Sigma_1^1$  definition.  $\square$

### 3 Ingredient 1: coding $\Delta_1^1$ functions

The proof of Theorem 1.1 involves several technical methods of rather general nature, which we present in the three following sections.

Recall that  $\omega_1^{\text{CK}}$  is the least non-recursive (= the least non- $\Delta_1^1$ ) ordinal. If  $\alpha < \omega_1^{\text{CK}}$  then let  $\mathcal{F}_\alpha$  be the set of all  $\Delta_1^1$  LR order preserving maps  $F : \langle \omega^\omega; \preceq \rangle \rightarrow \langle 2^\alpha; \leq_{1\text{ex}} \rangle$ , so that

$$x \preceq y \implies F(x) \leq_{1\text{ex}} F(y) \quad \text{for all } x, y \in \omega^\omega. \quad (1)$$

Such a function  $F$  has to be  $\approx$ -invariant on  $\omega^\omega$ . Let  $\mathcal{F} = \bigcup_{\alpha < \omega_1^{\text{CK}}} \mathcal{F}_\alpha$ .

If, in addition,  $X \subseteq \omega^\omega$  is a  $\Sigma_1^1$  set then let  $\mathcal{F}_X$  consist of all  $\Delta_1^1$  functions  $F \in \mathcal{F}$  such that

$$x, y \in X \text{ are } \preceq\text{-incomparable} \implies F(x) = F(y). \quad (2)$$

or equivalently,  $F(x) <_{1\text{ex}} F(y) \implies x \prec y$  for all  $x, y \in X$ .

Note that a function  $F \in \mathcal{F}_X$  has to be not just  $\approx$ -invariant, but also invariant w. r. t. the common equivalence hull of the relation  $\approx$  and the (non-equivalence) relation of being  $\preceq$ -incomparable. In particular, if for any  $x, y \in X$  there is  $z \in X$   $\preceq$ -incomparable with both  $x$  and  $y$ , then the only maps in  $\mathcal{F}_X$  are those constant on  $X$ .

**Definition 3.1.** Let, for  $x, y \in \omega^\omega$ :  $x \mathbf{E}_{\mathcal{F}} y$  iff  $\forall F \in \mathcal{F} (F(x) = F(y))$ ,

$$x \mathbf{E}_{\mathcal{F}_X} y \quad \text{iff} \quad \forall F \in \mathcal{F}_X (F(x) = F(y)) \quad (\text{here } X \subseteq \omega^\omega \text{ is } \Sigma_1^1). \quad \square$$

**Lemma 3.2.**  $E_{\mathcal{F}}$  is a smooth  $\Sigma_1^1$  equivalence relation, and if  $R(x, y)$  is a  $\Pi_1^1$  relation and  $\forall x, y (x E_{\mathcal{F}} y \implies R(x, y))$  then there is a single function  $F \in \mathcal{F}$  such that  $\forall x, y (F(x) = F(y) \implies R(x, y))$ .

Similarly, if  $X \subseteq \omega^\omega$  is a  $\Sigma_1^1$  set then  $E_{\mathcal{F}_X}$  is a smooth  $\Sigma_1^1$  equivalence relation, and if  $R(x, y)$  is a  $\Pi_1^1$  relation and  $\forall x, y (x E_{\mathcal{F}_X} y \implies R(x, y))$  then there is a function  $F \in \mathcal{F}_X$  such that  $\forall x, y (F(x) = F(y) \implies R(x, y))$ .

**Proof.** We concentrate on the first part; the result for the second part is pretty similar. We'll make use of an appropriate coding of functions in  $\mathcal{F}$ , based on a standard coding system of  $\Delta_1^1$  sets. A *code* will be a such-and-such pair  $f = \langle \varepsilon, k \rangle \in \omega^\omega \times \omega$ . We require that:

(I) the relation  $\leq_\varepsilon = \{\langle i, j \rangle : \varepsilon(2^i \cdot 3^j) = 0\}$  is a (non-strict) wellordering of the set  $\text{dom}(\leq_\varepsilon)$  — in this case, we let:

- $|\varepsilon| = \text{otp}(\varepsilon) < \omega_1$  be the order type of  $\leq_\varepsilon$ ,
- $\beta_\varepsilon : \text{dom}(\leq_\varepsilon) \xrightarrow{\text{onto}} |\varepsilon|$  be the order-preserving bijection,
- $H_\varepsilon : \omega^\omega \xrightarrow{\text{onto}} (\omega^\omega)^{|\varepsilon|}$  be the induced homeomorphism;

(II)  $k$  belongs to the set  $\mathbb{B} \subseteq \omega$  of *codes of  $\Delta_1^1$  sets*  $B \subseteq \omega^\omega \times \omega^\omega$ , so that it is assumed that  $\mathbb{B}$  is a  $\Pi_1^1$  set, and for any  $k \in \mathbb{B}$  a  $\Delta_1^1$  set  $B_k \subseteq \omega^\omega \times \omega^\omega$  is defined, and conversely, for any  $\Delta_1^1$  set  $B \subseteq \omega^\omega \times \omega^\omega$  there is a code  $k \in \mathbb{B}$  with  $B = B_k$ , and finally there exist two  $\Pi_1^1$  sets  $W, W' \subseteq \omega \times \omega^\omega \times \omega^\omega$  such that if  $k \in \mathbb{B}$  and  $x, y \in \omega^\omega$  then

$$\langle x, y \rangle \in B_k \iff W(k, x, y) \iff \neg W'(k, x, y);$$

(III) we define  $\mathbb{F}\text{un} = \{k \in \mathbb{B} : B_k \text{ is a total map } \omega^\omega \rightarrow \omega^\omega\}$ , the set of codes of all  $\Delta_1^1$  functions  $F : \omega^\omega \rightarrow \omega^\omega$  — this is still a  $\Pi_1^1$  set because the key condition  $\text{dom} B_f = \omega^\omega$  can be expressed by

$$\forall x \exists y \in \Delta_1^1(x) W(k, x, y),$$

where the quantifier  $\exists y \in \Delta_1^1(x)$  is known to preserve the type  $\Pi_1^1$ ;

(IV) if  $\varepsilon$  satisfies (I) and  $k \in \mathbb{F}\text{un}$  then let  $F_k^\varepsilon$  be the  $\Delta_1^1$  map  $\omega^\omega \rightarrow (\omega^\omega)^{|\varepsilon|}$  defined by  $F_k^\varepsilon(x) = H_\varepsilon(B_k(x))$  for all  $x \in \omega^\omega$ ;

**Definition 3.3.** Let  $\mathbb{F}$  be the set of all pairs  $f = \langle \varepsilon, k \rangle$  such that  $\varepsilon \in \omega^\omega$  satisfies (I),  $\varepsilon \in \Delta_1^1$ ,  $k \in \mathbb{F}\text{un}$ , and  $F_k^\varepsilon \in \mathcal{F}_{|\varepsilon|}$ .

If  $X \subseteq \omega^\omega$  is a  $\Sigma_1^1$  set then let  $\mathbb{F}(X) = \{\langle \varepsilon, k \rangle \in \mathbb{F} : F_k^\varepsilon \in \mathcal{F}_{|\varepsilon|}(X)\}$ .  $\square$

The following is a routine fact.

**Claim 3.4.**  $\mathbb{F} \subseteq \omega^\omega \times \omega$  is a countable  $\Pi_1^1$  set of  $\Delta_1^1$  elements, and  $\mathcal{F} = \{F_k^\varepsilon : \langle \varepsilon, k \rangle \in \mathbb{F}\}$ . If  $X \subseteq \omega^\omega$  is a  $\Sigma_1^1$  set then  $\mathbb{F}(X) \subseteq \omega^\omega \times \omega$  is a countable  $\Pi_1^1$  set of  $\Delta_1^1$  elements, and  $\mathcal{F}_X = \{F_k^\varepsilon : \langle \varepsilon, k \rangle \in \mathbb{F}(X)\}$ .  $\square$

In continuation of the proof of Lemma 3.2, note that

$$x \mathbf{E}_{\mathcal{F}} y \iff \forall \langle \varepsilon, k \rangle \in \mathbb{F} (F_k^\varepsilon(x) = F_k^\varepsilon(y)) \iff \forall \langle \varepsilon, k \rangle \in \mathbb{F} (B_k(x) = B_k(y)),$$

and this easily implies that  $\mathbf{E}_{\mathcal{F}}$  is  $\Sigma_1^1$  by Claim 3.4. Now prove the claim of Lemma 3.2 related to  $R(x, y)$ . We re-write the assumption as follows:

$$\forall x, y (\neg R(x, y) \implies \neg (x \mathbf{E}_{\mathcal{F}} y)),$$

or, equivalently by Claim 3.4, as

$$\forall x, y (\neg R(x, y) \implies \exists \langle \varepsilon, k \rangle \in \Delta_1^1 \underbrace{(\langle \varepsilon, k \rangle \in \mathbb{F} \wedge F_k^\varepsilon(x) \neq F_k^\varepsilon(y))}_{P(x, y; \varepsilon, k)}).$$

The relation  $P$  is expressible by a  $\Pi_1^1$  formula by means of (II) and Claim 3.4. It follows by Lemma 2.1 that there is a  $\Delta_1^1$  set  $W \subseteq \omega^\omega \times \omega^\omega$  satisfying  $\neg R(x, y) \implies W(x, y)$ , and a  $\Delta_1^1$  map  $\Phi(x, y) = \langle \varepsilon(x, y), k(x, y) \rangle : W \rightarrow \mathbb{F}$  such that we have  $F_{k(x, y)}^{\varepsilon(x, y)}(x) \neq F_{k(x, y)}^{\varepsilon(x, y)}(y)$  for all  $\langle x, y \rangle \in W$  — then, in particular, for all  $x, y$  with  $\neg R(x, y)$ .

The range  $H = \{\Phi(x, y) : \langle x, y \rangle \in W\}$  is then a  $\Sigma_1^1$  subset of the (countable)  $\Pi_1^1$  set  $\mathbb{F}$ . By Separation, there is a  $\Delta_1^1$  set  $D$  with  $H \subseteq D \subseteq \mathbb{F}$ . As a countable  $\Delta_1^1$  set, it admits a  $\Delta_1^1$  enumeration  $D = \{\langle \varepsilon_n, k_n \rangle : n \in \mathbb{N}\}$ , and by construction  $\forall n (F_{k_n}^{\varepsilon_n}(x) = F_{k_n}^{\varepsilon_n}(y))$  implies  $R(x, y)$ . Let

$$F(x) = F_{k_0}^{\varepsilon_0}(x) \wedge F_{k_1}^{\varepsilon_1}(x) \wedge F_{k_2}^{\varepsilon_2}(x) \wedge \dots$$

for  $x \in \omega^\omega$ . Then  $F \in \mathcal{F}$  and  $F(x) = F(y) \implies R(x, y)$ .  $\square$  (Lemma 3.2)

## 4 Ingredient 2: invariant separation

In the assumptions of Theorem 1.1, let  $\mathbf{E}$  be a  $\Sigma_1^1$  equivalence relation containing  $\approx$  (so that  $x \approx y$  implies  $x \mathbf{E} y$ ). A set  $X \subseteq \omega^\omega$  is *downwards  $\preceq$ -closed in each  $\mathbf{E}$ -class* iff we have  $x \in X \implies y \in X$  whenever  $x \mathbf{E} y$  and  $y \preceq x$ . The notion of a set *upwards  $\preceq$ -closed in each  $\mathbf{E}$ -class* is similar.

**Lemma 4.1.** *Let  $E$  be a  $\Sigma_1^1$  equivalence relation containing  $\approx$ ,  $X, Y$  be disjoint  $\Sigma_1^1$  sets, satisfying  $y \not\prec x$  whenever  $x \in X \wedge y \in Y \wedge x E y$ . Then there is a  $\Delta_1^1$  set  $Z$ , downwards  $\prec$ -closed in each  $E$ -class and satisfying  $X \subseteq Z$  and  $Y \cap Z = \emptyset$ .*

**Proof.** Let  $Y' = \{y' : \exists y \in Y (y \prec y')\}$ ; still  $Y' \cap X = \emptyset$  and  $Y'$  is  $\Sigma_1^1$ . Using Separation, define an increasing sequence of sets

$$X = X_0 \subseteq A_0 \subseteq X_1 \subseteq A_1 \subseteq \dots \subseteq X_n \subseteq A_n \subseteq \dots \subseteq \omega^\omega \setminus Y'$$

so that  $A_n \in \Delta_1^1$  and  $X_{n+1} = \{x' \in \omega^\omega : \exists x \in A_n (x' E x \wedge x' \prec x)\}$  for all  $n$ . If  $A_n \cap Y' = \emptyset$  then  $X_{n+1} \cap Y' = \emptyset$  as well since  $Y'$  is upwards closed, which justifies the inductive construction. Furthermore, a proper execution of the construction yields the final set  $Z = \bigcup_n A_n = \bigcup_n X_n$  in  $\Delta_1^1$ . (We refer to the proof of an “invariant” effective separation theorem in [1] or a similar construction in [4, Lemma 10.4.2].) Note that by construction  $X \subseteq Z$ , but  $Z \cap Y = \emptyset$ , and  $Z$  is downwards  $\prec$ -closed in each  $E$ -class.  $\square$

**Corollary 4.2.** *Let  $E$  be  $E_{\mathcal{F}}$ . If  $X, Y \subseteq \omega^\omega$  are disjoint  $\Sigma_1^1$  sets and  $[X]_E \cap [Y]_E \neq \emptyset$  then there are points  $x \in X, y \in Y$  with  $x E y$  and  $y \prec x$ .*

**Proof.** Otherwise by Lemma 4.1 there is a  $\Delta_1^1$  set  $Z$  such that  $X \subseteq Z$  and  $Y \cap Z = \emptyset$ , and downwards  $\prec$ -closed in each  $E$ -class. Then, by Lemma 3.2, there is a function  $F \in \mathcal{F}$  such that  $x \in Z \implies y \in Z$  holds whenever  $F(x) = F(y)$  and  $x \prec y$ . It follows that the derived function

$$G(x) = \begin{cases} F(x)^{\wedge 0}, & \text{whenever } x \in Z \\ F(x)^{\wedge 1}, & \text{whenever } x \in \omega^\omega \setminus Z \end{cases}$$

belongs to  $\mathcal{F}$ . Thus if  $x \in Z$  and  $y \notin Z$ , say  $x \in X$  and  $y \in Y$ , then  $G(x) \neq G(y)$  and hence  $x \not E y$ , a contradiction.  $\square$

## 5 Ingredient 3: the Gandy – Harrington forcing

The Gandy – Harrington forcing notion  $\mathbb{P}$  is the set of all  $\Sigma_1^1$  sets  $\emptyset \neq X \subseteq \omega^\omega$ , ordered so that smaller sets are stronger conditions. We also define  $\mathbb{P}_n$  ( $n \geq 2$ ) to be the set of all  $\Sigma_1^1$  sets  $\emptyset \neq X \subseteq (\omega^\omega)^n$ .

It is known that  $\mathbb{P}$  adds a point of  $\omega^\omega$ , whose name will be  $\dot{x}$ .

Together with  $\mathbb{P}$ , some other related forcing notions will be considered below, for instance, the product  $\mathbb{P}^2 = \mathbb{P} \times \mathbb{P}$  which consists of all cartesian

products of the form  $X \times Y$ , where  $X, Y \in \mathbb{P}$ . It follows from the above that  $\mathbb{P}^2$  forces a pair of points of  $2^\omega$ , whose name will be  $\langle \dot{x}_{1e}, \dot{x}_{ri} \rangle$ .

There is another important subforcing introduced in [2]. If  $\mathbb{E}$  is a  $\Sigma_1^1$  equivalence relation on  $\omega^\omega$  then let  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$  consist of all sets of the form  $X \times Y$ , where  $X, Y \in \mathbb{P}$  and  $(X \times Y) \cap \mathbb{E} \neq \emptyset$ .

A condition  $X \times Y$  in  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$  is *saturated* iff  $[X]_{\mathbb{E}} = [Y]_{\mathbb{E}}$ .

**Lemma 5.1.** *If  $X \times Y$  is a condition in  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$  then there is a stronger saturated subcondition  $X' \times Y'$  in  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$ .*

**Proof.**  $X' = \{x \in X : \exists y \in Y (x \mathbb{E} y)\}$ ,  $Y' = \{y \in Y : \exists x \in X (x \mathbb{E} y)\}$ .  $\square$

**Remark 5.2.** If  $X \times Y$  is a saturated condition in  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$ , and  $\emptyset \neq X' \subseteq X$  is a  $\Sigma_1^1$  set, then  $Y' = Y \cap [X']_{\mathbb{E}}$  is  $\Sigma_1^1$  and  $X' \times Y'$  is still a saturated condition in  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$ . It follows that  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$  forces a pair of  $\mathbb{P}$ -generic reals, whose names will be  $\dot{x}_{1e}$  and  $\dot{x}_{ri}$  as above.

**Lemma 5.3** (2.9 in [2]). *Suppose that  $\mathbb{E}$  is a smooth  $\Sigma_1^1$  equivalence relation. Then  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$  forces  $\dot{x}_{1e} \mathbb{E} \dot{x}_{ri}$ .*  $\square$

Note that Lemma 5.3, generally speaking, fails in the non-smooth case.

The next result will be pretty important.

**Lemma 5.4** (2.9 in [2]). *Suppose that  $\preceq$  is a  $\Delta_1^1$  PQO on  $\omega^\omega$ , and for any  $A \in \mathbb{P}$  there is a  $\Sigma_1^1$  equivalence relation  $\mathbb{E}_A$  on  $\omega^\omega$  such that if  $A \subseteq B$  then  $x \mathbb{E}_A y$  implies  $x \mathbb{E}_B y$ . Assume that  $X^* \in \mathbb{P}$ , and if  $B \in \mathbb{P}$ ,  $B \subseteq X^*$  then  $B \times B$  does **not** ( $\mathbb{P} \times_{\mathbb{E}_B} \mathbb{P}$ )-force that  $\dot{x}_{1e}, \dot{x}_{ri}$  are  $\preceq$ -comparable.*

*Then  $X^*$  is not  $\preceq$ -thin, in other words, there is a perfect set  $Y \subseteq X^*$  of pairwise  $\preceq$ -incomparable elements.*  $\square$

The forcing  $\mathbb{P}$ , as well as some of its derivatives like  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$ , will be used below as forcing notions over the ground set universe  $\mathbf{V}$ .

**Lemma 5.5** (see [1, 2]). *If  $X \in \mathbb{P}$  then  $X$   $\mathbb{P}$ -forces that  $\dot{x} \in X$ . Moreover if  $\Phi(x)$  is a  $\Pi_2^1$  formula and  $\Phi(x)$  holds for all  $x \in X$  then  $X$   $\mathbb{P}$ -forces that  $\dot{x}$  satisfies  $\Phi(X)$ .*

*The same is true for other similar forcing notions like  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$ .*  $\square$

Here (and below in some cases), given a  $\Sigma_1^1$  (or  $\Pi_1^1$ ) set  $X$  in the ground universe  $\mathbf{V}$ , **we denote by the same letter  $X$  the extended set** (i. e., defined by the same formula) **in any generic extension of  $\mathbf{V}$** . By the Shoenfield absoluteness theorem, there is no ambiguity here. See [5, 2.4] in more detail.

## 6 Bounding thin partial orderings

Here we prove claim (i) of Theorem 1.1. We'll make use of the family  $\mathcal{F}$  of  $\Delta_1^1$  functions, introduced in Section 3, and the corresponding smooth  $\Sigma_1^1$  equivalence relation  $\mathbf{E} = \mathbf{E}_{\mathcal{F}}$ . Then  $\approx$  is a subrelation of  $\mathbf{E}$  by Lemma 3.2.

The following partition on cases is quite common in this sort of proofs.

**Case 1:**  $\approx$  and  $\mathbf{E}$  coincide on  $X^*$ , so that  $x \mathbf{E} y \implies x \approx y$  for  $x, y \in X^*$ . Then, by Lemma 3.2, there is a single function  $F \in \mathcal{F}$  such that  $F(x) = F(y)$  implies  $x \approx y$  for all  $x, y \in X^*$ , as required.

**Case 2:**  $\approx$  is a *proper* subrelation of  $\mathbf{E}$  on  $X^*$ , hence, the  $\Sigma_1^1$  set

$$V^* = \{x \in X^* : \exists y \in X^* (x \not\approx y \wedge x \mathbf{E} y)\}$$

is non-empty. Our final goal will be to infer a contradiction; then the result for Case 1 proves Claim (i) of the theorem.

Note that  $V^* \times V^*$  is a saturated condition in  $\mathbb{P} \times_{\mathbf{E}} \mathbb{P}$ .

**Lemma 6.1.** *Condition  $V^* \times V^*$  ( $\mathbb{P} \times_{\mathbf{E}} \mathbb{P}$ )-forces that  $\dot{x}_{1e}$  and  $\dot{x}_{ri}$  are  $\preceq$ -incomparable.*

**Proof.** Suppose to the contrary that a subcondition  $Y \times Z$  either forces  $\dot{x}_{1e} \approx \dot{x}_{ri}$  or forces  $\dot{x}_{1e} \prec \dot{x}_{ri}$ . We will get a contradiction in both cases. Note that  $Y, Z \subseteq V^*$  are non-empty  $\Sigma_1^1$  sets and  $[Y]_{\mathbf{E}} \cap [Z]_{\mathbf{E}} \neq \emptyset$ .

**Case A:**  $Y \times Z$  forces  $\dot{x}_{1e} \approx \dot{x}_{ri}$ .

**Subcase A1:** the  $\Sigma_1^1$  set  $W = \{(y, y') \in Y \times Y : y \mathbf{E} y' \wedge y' \not\approx y\}$  is empty, or in other words  $\mathbf{E}$  coincides with  $\approx$  on  $Y$ . By the non-emptiness of  $V^*$  at least one of the  $\Sigma_1^1$  sets

$$B = \{x : \exists y \in Y (x \mathbf{E} y \wedge x \not\approx y)\}, \quad B' = \{x : \exists y \in Y (x \mathbf{E} y \wedge y \not\approx x)\}$$

is non-empty; assume that, say,  $B \neq \emptyset$ . Consider the  $\Sigma_1^1$  set

$$A = \{x : \exists y \in Y (x \mathbf{E} y \wedge x \preceq y)\}; \quad Y \subseteq A.$$

Then  $A \cap B = \emptyset$ ,  $A$  is downwards closed while  $B$  is upwards closed in each  $\mathbf{E}$ -class, therefore  $y \not\approx x$  whenever  $x \in A$ ,  $y \in B$ , and  $x \mathbf{E} y$ . Then  $[A]_{\mathbf{E}} \cap [B]_{\mathbf{E}} = \emptyset$  by Corollary 4.2. Yet by definition  $[Y]_{\mathbf{E}} \cap [B]_{\mathbf{E}} \neq \emptyset$  and  $Y \subseteq A$ , which is a contradiction.

**Subcase A2:**  $W \neq \emptyset$ . Then the forcing notion  $\mathbb{P}(W)$  of all non-empty  $\Sigma_1^1$  sets  $P \subseteq W$  adds pairs of  $\mathbb{P}$ -generic (separately) reals  $y, y' \in Y$  which belong to  $W$  and satisfy  $y' \mathbf{E} y$  and  $y' \not\approx y$ , by Lemma 5.5.



If  $P \in \mathbb{P}(W)$  then obviously  $[\text{dom } P]_{\mathbb{E}} = [\text{ran } P]_{\mathbb{E}}$ .

Consider a more complex forcing notion  $\mathcal{P} = \mathbb{P}(W) \times_{\mathbb{E}} \mathbb{P}$  of all pairs  $P \times Z'$ , where  $P \in \mathbb{P}(W)$ ,  $Z' \in \mathbb{P}$ ,  $Z' \subseteq Z$ , and  $[\text{dom } P]_{\mathbb{E}} \cap [Z']_{\mathbb{E}} \neq \emptyset$ . For instance,  $W \times Z \in \mathbb{P}(W) \times_{\mathbb{E}} \mathbb{P}$ . Then  $\mathcal{P}$  adds a pair  $\langle \dot{x}_{1e}, \dot{x}_{ri} \rangle \in W$  and a separate real  $\dot{x} \in B$  such that both pairs  $\langle \dot{x}_{1e}, \dot{x} \rangle$  and  $\langle \dot{x}_{ri}, \dot{x} \rangle$  are  $(\mathbb{P} \times_{\mathbb{E}} \mathbb{P})$ -generic, hence, we have  $\dot{x}_{1e} \approx \dot{x} \approx \dot{x}_{ri}$  (in the extended universe  $\mathbf{V}[\dot{x}_{1e}, \dot{x}_{ri}, \dot{x}]$ ) by the choice of  $Y \times Z$ . On the other hand,  $\dot{x}_{1e} \not\approx \dot{x}_{ri}$  by Lemma 5.5, since the pair belongs to  $W$ , which is a contradiction.

**Case B:**  $Y \times Z$  forces  $\dot{x}_{1e} \prec \dot{x}_{ri}$ .

**Subcase B1:** the  $\Sigma_1^1$  set  $W = \{\langle y, z \rangle \in Y \times Z : z \mathbb{E} y \wedge z \preceq y\}$  is empty. Then the  $\Sigma_1^1$  sets

$$Y_0 = \{y' : \exists y \in Y (y \mathbb{E} y' \wedge y' \preceq y)\}, \quad Z_0 = \{z' : \exists z \in Z (z \mathbb{E} z' \wedge z \preceq z')\}$$

are disjoint and  $\preceq$ -closed resp. downwards and upwards, hence we have  $[Z_0]_{\mathbb{E}} \cap [Y_0]_{\mathbb{E}} = \emptyset$  by Corollary 4.2. However  $[Z]_{\mathbb{E}} \cap [Y]_{\mathbb{E}} \neq \emptyset$ , which is a contradiction as  $Z \subseteq Z_0$ ,  $Y \subseteq Y_0$ .

**Subcase B2:**  $W \neq \emptyset$ . Consider the forcing  $\mathbb{P}(W)$  of all non-empty  $\Sigma_1^1$  sets  $P \subseteq W$ ; if  $P \in \mathbb{P}(W)$  then obviously  $[\text{dom } P]_{\mathbb{E}} = [\text{ran } P]_{\mathbb{E}}$ . Consider a more complicated forcing  $\mathbb{P}(W) \times_{\mathbb{E}} \mathbb{P}(W)$  of all products  $P \times Q$ , where  $P, Q \in \mathbb{P}(W)$  and  $[\text{dom } P]_{\mathbb{E}} \cap [\text{dom } Q]_{\mathbb{E}} \neq \emptyset$ . In particular  $W \times W \in \mathbb{P}(W) \times_{\mathbb{E}} \mathbb{P}(W)$ .

Let  $\langle x, y, x', y' \rangle$  be a  $\mathbb{P}(W) \times_{\mathbb{E}} \mathbb{P}(W)$ -generic quadruple in  $W \times W$ , so that both  $\langle x, y \rangle \in W$  and  $\langle x', y' \rangle \in W$  are  $\mathbb{P}(W)$ -generic pairs in  $W$ , and both  $y \preceq x$  and  $y' \preceq x'$  hold by the definition of  $W$ . On the other hand, an easy argument shows that both criss-cross pairs  $\langle x, y' \rangle \in X \times Y$  and  $\langle x', y \rangle \in X \times Y$  are  $\mathbb{P} \times_{\mathbb{E}} \mathbb{P}$ -generic, hence  $x \prec y'$  and  $x' \prec y$  by the choice of  $X \times Y$ . Altogether  $y \preceq x \prec y' \preceq x' \prec y$ , which is a contradiction.  $\square$

To accomplish the proof of (i) of Theorem 1.1, note that by Lemma 6.1 and Lemma 5.4 (with  $\mathbb{E}_A = \mathbb{E}$  for all  $A$ ) there is a perfect 2wise  $\approx$ -inequivalent set, so  $\preceq$  is not thin, contrary to our assumptions.

## 7 Decomposing thin partial orderings

We prove **claim (ii) of Theorem 1.1** in this Section. Let  $U^*$  be the  $\Sigma_1^1$  set of all reals  $x \in X^*$  such that there is no  $\Delta_1^1$   $\preceq$ -chain  $C$  containing  $x$ .

**We assume to the contrary that  $U^* \neq \emptyset$ .**

The proof will make heavy use of the functions in families of the form  $\mathcal{F}_X$ , introduced in Section 3. If  $X \subseteq \omega^\omega$  is a  $\Sigma_1^1$  set then  $\mathbb{E}_X = \mathbb{E}_{\mathcal{F}_X}$  is a smooth  $\Sigma_1^1$  equivalence relation by Lemma 3.2.

If  $X \subseteq X'$  then  $\mathcal{F}'_X \subseteq \mathcal{F}_X$ , and hence  $x \mathbf{E}_X y$  implies  $x \mathbf{E}_{X'} y$ .

**Corollary 7.1** (of Lemma 5.3). *If  $X \subseteq U^*$  is a non-empty  $\Sigma_1^1$  set then the condition  $X \times X$  ( $\mathbb{P} \times_{\mathbf{E}_X} \mathbb{P}$ )-forces that  $\dot{\mathbf{x}}_{1e} \mathbf{E}_X \dot{\mathbf{x}}_{ri}$ .  $\square$*

**Lemma 7.2.** *Let  $X \subseteq U^*$  be a non-empty  $\Sigma_1^1$  set. Then  $X \times X$  does **not** ( $\mathbb{P} \times_{\mathbf{E}_X} \mathbb{P}$ )-force that  $\dot{\mathbf{x}}_{1e}, \dot{\mathbf{x}}_{ri}$  are  $\preceq$ -comparable.*

**Proof.** Suppose to the contrary that  $X \times X$  forces the comparability. Then there is a subcondition  $Y \times Z$  which either forces  $\dot{\mathbf{x}}_{1e} \approx \dot{\mathbf{x}}_{ri}$  or forces  $\dot{\mathbf{x}}_{1e} \prec \dot{\mathbf{x}}_{ri}$ ;  $Y, Z \subseteq X$  are non-empty  $\Sigma_1^1$  sets and  $[Y]_{\mathbf{E}_X} \cap [Z]_{\mathbf{E}_X} \neq \emptyset$ .

**Case A:**  $Y \times Z$  forces  $\dot{\mathbf{x}}_{1e} \approx \dot{\mathbf{x}}_{ri}$ .

**Subcase A1:** the  $\Sigma_1^1$  set  $W = \{\langle y, y' \rangle \in Y \times Y : y \mathbf{E}_X y' \wedge y' \not\approx y\}$  is empty. Then  $Y$  is a  $\preceq$ -chain: indeed if  $x, y \in Y$  are  $\preceq$ -incomparable then by definition we have  $x \mathbf{E}_X y$ , hence  $x \approx y$ , contradiction. Let  $C$  be the  $\Pi_1^1$  set of all reals  $\preceq$ -comparable with each  $y \in Y$ ; then  $Y \subseteq C$ . By Separation there is a  $\Delta_1^1$  set  $D$ ,  $Y \subseteq D \subseteq C$ . Let  $C'$  be the  $\Pi_1^1$  set of all reals in  $D$ ,  $\preceq$ -comparable with each  $d \in D$ ; then  $Y \subseteq C' \subseteq D$ . Take any  $\Delta_1^1$  set  $B$  with  $Y \subseteq B \subseteq C'$ . By construction  $B$  is a  $\Delta_1^1$   $\preceq$ -chain with  $\emptyset \neq Y \subseteq B$ , contrary to the definition of  $U^*$ .

**Subcase A2:**  $W \neq \emptyset$ : yields a contradiction similarly to Subcase A2 in the proof of Lemma 6.1.

**Case B:**  $Y \times Z$  forces  $\dot{\mathbf{x}}_{1e} \prec \dot{\mathbf{x}}_{ri}$ .

**Subcase B1:** the  $\Sigma_1^1$  set  $W = \{\langle y, z \rangle \in Y \times Z : y \mathbf{E}_X z \wedge y \not\prec z\}$  is non-empty. Let  $Y' = \text{dom } W$ . As  $Y' \subseteq X$ , the condition  $Y' \times Y'$  ( $\mathbb{P} \times_{\mathbf{E}_X} \mathbb{P}$ )-forces that  $\dot{\mathbf{x}}_{1e}, \dot{\mathbf{x}}_{ri}$  are  $\preceq$ -comparable. Therefore by the result in Case A there is a condition  $A \times B$  in  $\mathbb{P} \times_{\mathbf{E}_X} \mathbb{P}$ , with  $A \cup B \subseteq Y'$ , which forces  $\dot{\mathbf{x}}_{1e} \prec \dot{\mathbf{x}}_{ri}$ ; for if it forces  $\dot{\mathbf{x}}_{ri} \prec \dot{\mathbf{x}}_{1e}$  then just consider  $B \times A$  instead of  $A \times B$ . Consider the forcing notion  $\mathcal{P}$  of all non- $\emptyset$   $\Sigma_1^1$  sets of the form  $P \times B'$ , where

$$P \subseteq W, \quad \text{dom } P \subseteq A, \quad B' \subseteq B, \quad \text{and} \quad [B']_{\mathbf{E}_X} = [\text{dom } P]_{\mathbf{E}_X} = [\text{ran } P]_{\mathbf{E}_X}.$$

For instance if  $B' = B$  and  $P = \{\langle x, y \rangle \in W : x \in A\}$  then  $P \times B' \in \mathcal{P}$ . Then  $\mathcal{P}$  forces a pair  $\langle \dot{\mathbf{x}}_{1e}, \dot{\mathbf{x}}_{ri} \rangle \in W$  and a separate real  $\dot{\mathbf{x}} \in B$  such that both pairs  $\langle \dot{\mathbf{x}}_{1e}, \dot{\mathbf{x}} \rangle$  and  $\langle \dot{\mathbf{x}}_{ri}, \dot{\mathbf{x}} \rangle$  are ( $\mathbb{P} \times_{\mathbf{E}_X} \mathbb{P}$ )-generic. It follows that  $\mathcal{P}$  forces both  $\dot{\mathbf{x}}_{1e} \prec \dot{\mathbf{x}}$  (as this pair belongs to  $A \times B$ ) and  $\dot{\mathbf{x}} \prec \dot{\mathbf{x}}_{ri}$  (as this pair belongs to  $Y \times Z$ ), hence, forces  $\dot{\mathbf{x}}_{1e} \prec \dot{\mathbf{x}}_{ri}$ . On the other hand  $\mathcal{P}$  forces  $\dot{\mathbf{x}}_{1e} \not\prec \dot{\mathbf{x}}_{ri}$  (as this pair belongs to  $W$ ), which is a contradiction.

**Subcase B2:**  $W = \emptyset$ , in other words, if  $y \in Y$ ,  $z \in Z$ , and  $y \mathbf{E}_X z$  then  $y \prec z$  strictly. Then by Lemma 4.1 there is a  $\Delta_1^1$  set  $C \subseteq \omega^\omega$ , downwards

$\preceq$ -closed in each  $E_X$ -class, such that  $Y \subseteq C$  and still  $Z \cap C = \emptyset$ . We claim that, moreover,

*if  $y \in C \cap X$ ,  $z \in X \setminus C$ , and  $y E_X z$ , then  $y \prec z$ .*

Indeed otherwise, the following  $\Sigma_1^1$  set

$$H_0 = \{z \in X \setminus C : \exists y \in C \cap X (y E_X z \wedge y \not\prec z)\} \subseteq X$$

is non- $\emptyset$ . As above, there is a saturated condition  $H \times H'$  in  $\mathbb{P} \times_{E_X} \mathbb{P}$ , with  $H \cup H' \subseteq H_0$ , which forces  $\dot{x}_{1e} \prec \dot{x}_{ri}$ , and then  $z \prec z'$  holds whenever  $\langle z, z' \rangle \in H \times H'$  and  $z E_X z'$ . By construction the  $\Sigma_1^1$  set

$$C_1 = \{y \in C \cap X : \exists z' \in H' (y E_X z' \wedge y \not\prec z')\}$$

satisfies  $[C_1]_{E_X} = [H]_{E_X} = [H']_{E_X}$ , hence  $C_1 \times H$  is a condition in  $\mathbb{P} \times_{E_X} \mathbb{P}$ . Let  $\langle y_1, z \rangle \in C_1 \times H$  be any  $(\mathbb{P} \times_{E_X} \mathbb{P})$ -generic pair. Then  $y_1 E_X z$  by Corollary 7.1, and, by the choice of  $X$  and the result in Case A, we have  $y_1 \prec z$  or  $z \prec y_1$ . However by construction  $y_1 \in C$ ,  $z \notin C$ , and  $C$  is downwards closed in each  $E_X$ -class. Thus in fact  $y_1 \prec z$ . Therefore, for all  $z' \in H'$ , if  $y_1 E_X z'$  then  $y_1 \prec z \prec z'$ , which contradicts to  $y_1 \in C_1$ .

Thus indeed  $y \prec z$  holds whenever  $y \in C \cap X$ ,  $z \in X \setminus C$ , and  $y E_X z$ . By Lemma 3.2 *there is a single function  $F \in \mathcal{F}_X$  such that if  $y \in C \cap X$ ,  $z \in X \setminus C$ , and  $F(y) = F(z)$ , then  $y \prec z$ .*

We claim that *the derived function*

$$G(x) = \begin{cases} F(x)^{\wedge 0}, & \text{whenever } x \in C \\ F(x)^{\wedge 1}, & \text{whenever } x \in \omega^\omega \setminus C \end{cases}$$

*belongs to  $\mathcal{F}_X$ .* First of all, still  $G \in \mathcal{F}$  since  $C$  is downwards  $\preceq$ -closed in each  $\mathbb{P} \times_{E_X} \mathbb{P}$ -class. Now suppose that  $z, y \in X$  and  $G(y) <_{1ex} G(z)$ . Then either  $F(y) <_{1ex} F(z)$ , or  $F(z) = F(y)$  and  $y \in C$  but  $z \notin C$ . In the “either” case immediately  $y \prec z$  since  $F \in \mathcal{F}_X$ <sup>2</sup>. In the “or” case we have  $y \prec z$  by the choice of  $F$  and the definition of  $G$ . Thus  $G \in \mathcal{F}_X$ .

Now pick any pair of reals  $y \in Y$  and  $z \in Z$  with  $y E_X z$ . Then we have  $G(x) = G(y)$  since  $G \in \mathcal{F}_X$ . But  $y \in C$  and  $z \notin C$  hold since  $Y \subseteq C$  and  $Z \cap C = \emptyset$  by construction, and in this case surely  $G(y) \neq G(z)$  by the definition of  $G$ . This contradiction completes the proof of Lemma 7.2.  $\square$

Lemma 7.2 plus Lemma 5.4 imply claim (ii) of Theorem 1.1.

$\square$  (Theorem 1.1)

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<sup>2</sup> The family  $\mathcal{F}$  would not work in the passage; here we have to use  $\mathcal{F}_X$  instead.

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