

The parameterfree Comprehension does not imply the full Comprehension in the 2nd order Peano arithmetic*

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Abstract

The parameter-free part \mathbf{PA}_2^* of \mathbf{PA}_2 , the 2nd order Peano arithmetic, is considered. We make use of a product/iterated Sacks forcing to define an ω -model of $\mathbf{PA}_2^* + \mathbf{CA}(\Sigma_2^1)$, in which an example of the full Comprehension schema \mathbf{CA} fails. Using Cohen's forcing, we also define an ω -model of \mathbf{PA}_2^* , in which not every set has its complement, and hence the full \mathbf{CA} fails in a rather elementary way.

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1 Introduction

Discussing the structure and deductive properties of the second order Peano arithmetic \mathbf{PA}_2 , Kreisel [9, § III, page 366] wrote that the selection of subsystems “is a central problem”. In particular, Kreisel notes, that

[...] if one is convinced of the significance of something like a given axiom schema, it is natural to study details, such as the effect of parameters.

Recall that *parameters* in this context are free variables in various axiom schemata in \mathbf{PA} , \mathbf{ZFC} , and other similar theories. Thus the most obvious way to study “the effect of parameters” is to compare the strength of a given axiom schema S with its parameter-free subschema S^* . (The asterisk will mean the parameter-free subschema in this paper.)

Some work in this direction was done in the early years of modern set theory. In particular Guzikki [6] proved that the Levy-style generic collapse (see, e.g., Levy [11] and Solovay [18]) of all cardinals $\omega_\alpha^{\mathbf{L}}$, $\alpha < \omega_1^{\mathbf{L}}$, results in a generic extension of \mathbf{L} in which the (countable) choice schema \mathbf{AC} , in the language of \mathbf{PA}_2 , fails but its parameter-free subschema \mathbf{AC}^* holds, so that \mathbf{AC}^* is strictly weaker than \mathbf{AC} . This can be compared with an opposite result for the *dependent choice* schema \mathbf{DC} , in the language of \mathbf{PA}_2 , which is equivalent to its parameter-free subschema \mathbf{DC}^* by a simple argument given in [6].

Some results related to parameter-free versions of the Separation and Replacement axiom schemata in \mathbf{ZFC} also are known from [3, 12, 14].

This paper is devoted to the role of parameters in the *comprehension schema* \mathbf{CA} of \mathbf{PA}_2 . Let \mathbf{PA}_2^* be the subtheory of \mathbf{PA}_2 in which the full schema \mathbf{CA} is replaced by its parameter-free version \mathbf{CA}^* , and the Induction principle is formulated as a schema rather than one sentence. The following Theorems 1.1 and 1.2 are our main results.

Theorem 1.1. *Let Cohen be the Cohen forcing for adding a generic subset of ω . Let Cohen^ω be the finite-support product. Suppose that $\langle x_i \rangle_{i < \omega_1}$ is a sequence Cohen^ω -generic over \mathbf{L} , the constructible universe.*

Let $X = (\mathcal{P}(\omega) \cap \mathbf{L}) \cup \{x_i : i < \omega\}$. Then $\langle \omega; X \rangle$ is a model of \mathbf{PA}_2^ , but not a model of \mathbf{CA} as X does not contain the complements $\omega \setminus x_i$.*

Thus \mathbf{CA} , even in the particular form claiming that every set has its complement, is not provable in \mathbf{PA}_2^ .*

It is quite obvious that a subtheory like \mathbf{PA}_2^* , that does not allow such a fundamental thing as the complement formation, is unacceptable. This is why we adjoin $\mathbf{CA}(\Sigma_2^1)$, i.e., the full \mathbf{CA} (with parameters) restricted to Σ_2^1 formulas, in the next theorem, to obtain a more plausible subsystem.

Theorem 1.2. *There is a generic extension $\mathbf{L}[G]$ of \mathbf{L} and a set $M \in \mathbf{L}[G]$, such that $\mathcal{P}(\omega) \cap \mathbf{L} \subseteq M \subseteq \mathcal{P}(\omega)$ and $\langle \omega; M \rangle$ is a model of $\mathbf{PA}_2^* + \mathbf{CA}(\Sigma_2^1)$ but not a model of \mathbf{PA}_2 .*

Therefore \mathbf{CA} is not provable even in $\mathbf{PA}_2^ + \mathbf{CA}(\Sigma_2^1)$.*

Theorem 1.2 will be established by means of a complex product/iteration of the Sacks forcing and the associated coding by degrees of constructibility, approximately as discussed in [13, page 143], around Theorem T3106.

Identifying the theories with their deductive closures, we may present the concluding statements of Theorems 1.1 and 1.2 as resp.

$$\mathbf{PA}_2^* \subsetneq \mathbf{PA}_2 \quad \text{and} \quad (\mathbf{PA}_2^* + \mathbf{CA}(\Sigma_2^1)) \subsetneq \mathbf{PA}_2.$$

Studies on subsystems of \mathbf{PA}_2 have discovered many cases in which $S \subsetneq S'$ holds for a given pair of subsystems S, S' , see e.g. [17]. And it is a rather typical case that such a strict extension is established by demonstrating that S' proves the consistency of S . One may ask whether this is the case for the results in the displayed line above. The answer is in the negative: namely the theories \mathbf{PA}_2^* , $\mathbf{PA}_2^* + \mathbf{CA}(\Sigma_2^1)$, and the full \mathbf{PA}_2 happen to be equiconsistent by a result in [4], also mentioned in [15]. This equiconsistency result also follows from a somewhat sharper theorem in [16, 1.5].¹

2 Preliminaries

Following [1, 9, 17] we define the second order Peano arithmetic \mathbf{PA}_2 as a theory in the language $\mathcal{L}(\mathbf{PA}_2)$ with two sorts of variables – for natural numbers and for sets of them. We use j, k, m, n for variables over ω and x, y, z for variables over $\mathcal{P}(\omega)$, reserving capital letters for subsets of $\mathcal{P}(\omega)$ and other sets. The axioms are as follows:

- (1) Peano's axioms for numbers.
- (2) The Induction schema $\Phi(0) \wedge \forall k (\Phi(k) \implies \Phi(k+1)) \implies \forall k \Phi(k)$, for every formula $\Phi(k)$ in $\mathcal{L}(\mathbf{PA}_2)$, and in $\Phi(k)$ we allow parameters, i.e., free variables other than k .²
- (3) Extensionality for sets.
- (4) The Comprehension schema \mathbf{CA} : $\exists x \forall k (k \in x \iff \Phi(k))$, for every formula Φ in which the variable x does not occur, and in Φ we allow parameters.

¹ The authors are thankful to Ali Enayat for the references to [4, 15, 16] in matters of this equiconsistency result.

² We cannot use Induction as one sentence because the Comprehension schema \mathbf{CA} is not assumed in full generality in the context of Theorem 1.1.

We let $\mathbf{CA}(\Sigma_2^1)$ be the full \mathbf{CA} restricted to Σ_2^1 formulas Φ .³

We let \mathbf{CA}^* be the parameter-free sub-schema of \mathbf{PA} (that is, $\Phi(k)$ contains no free variables other than k).

We let \mathbf{PA}_2^* be the subsystem of \mathbf{PA}_2 with \mathbf{CA} replaced by \mathbf{CA}^* .

Remark 2.1. In spite of Theorem 1.1, \mathbf{PA}_2^* proves \mathbf{CA} with parameters over ω (but not over $\mathcal{P}(\omega)$) allowed. Indeed suppose that Φ is $\Phi(k, m)$ in (4) and Φ has no other free variables. Arguing in \mathbf{PA}_2 , assume towards the contrary that the formula $\psi(m) := \exists x \forall k (k \in x \iff \Phi(k, m))$ holds not for all m . By Induction, take the least m for which $\psi(m)$ fails. This m is definable, and therefore it can be eliminated, and hence we have $\psi(m)$ for this m by \mathbf{CA}^* . This is a contradiction. \square

3 Extension by Cohen reals

Here we prove Theorem 1.1. We assume some knowledge of forcing and generic models, as e.g. in Kunen [10], especially Section IV.6 there on the “forcing over the universe” approach.

Recal that the Cohen forcing notion $\mathbf{Cohen} = 2^{<\omega}$ consists of all finite dyadic tuples including the empty tuple Λ . If $u, v \in 2^{<\omega}$ then $u \subset v$ means that v is a proper extension of u , whereas $u \subseteq v$ means $u \subset v \vee u = v$. The *finite-support product* $\mathbf{P} = (2^{<\omega})^\omega$ consists of all maps $p : \omega \rightarrow 2^{<\omega}$ such that $p(i) = \Lambda$ (the empty tuple) for all but finite $i < \omega$. The set \mathbf{P} is ordered opposite to the componentwise extension, so that $p \leq q$ (p is *stronger* as a forcing condition) iff $q(i) \subseteq p(i)$ for all $i < \omega$. The condition Λ^ω defined by $\Lambda^\omega(i) = \Lambda, \forall i$, is the \leq -largest (the weakest) element of \mathbf{P} .

We consider the set \mathbf{Perm} of all idempotent *permutations* of ω , that is, all bijections $\pi : \omega \xrightarrow{\text{onto}} \omega$ such that $\pi = \pi^{-1}$ and *the domain of nontriviality* $|\pi| = \{i : \pi(i) \neq i\}$ is finite. If $\pi \in \mathbf{Perm}$ and p is a function with $\text{dom } \pi = \omega$, then πp is defined by $\text{dom}(\pi p) = \omega$ and $(\pi p)(\pi(i)) = p(i)$ for all $i < \omega$, so formally $\pi p = p \circ \pi^{-1} = p \circ \pi$ (the superposition). In particular if $p \in \mathbf{P}$ then $\pi p \in \mathbf{P}$ and $|\pi p| = \pi''|p| = \{\pi(i) : i \in |p|\}$.

Proof (Theorem 1.1). We make use of Gödel’s *constructible universe* \mathbf{L} as *the ground model* for our forcing constructions. Suppose that $G \subseteq \mathbf{P}$ is a set \mathbf{P} -generic over \mathbf{L} . If $i < \omega$ then

- $G_i = \{p(i) : p \in G\} \subseteq 2^{<\omega}$ is a set $2^{<\omega}$ -generic (Cohen generic) over \mathbf{L} ,
- $a_i[G] = \bigcup G_i \in 2^\omega$ is a real Cohen generic over \mathbf{L} , and
- $x_i[G] = \{n : a_i(n) = 1\} \subseteq \omega$ is a subset of ω Cohen generic over \mathbf{L} .

³ A Σ_2^1 formula is any $\mathcal{L}(\mathbf{PA}_2)$ formula of the form $\forall x \exists y \Psi$, where Ψ does not contain quantified variables over $\mathcal{P}(\omega)$.

$$- X = X[G] = (\mathcal{P}(\omega) \cap \mathbf{L}) \cup \{x_i[G] : i < \omega\}.$$

Thus $X[G] \in \mathbf{L}[G]$ and $X[G]$ consists of all subsets of ω already in \mathbf{L} and all Cohen-generic sets $x_i[G]$, $i < \omega$.

We assert that the model $\langle \omega; X[G] \rangle$ proves Theorem 1.1.

The only thing to check is that $\langle \omega; X[G] \rangle$ satisfies \mathbf{CA}^* . For that purpose, assume that $\Phi(k)$ is a parameter-free $\mathcal{L}(\mathbf{PA}_2)$ formula with k the only free variable. Consider the set $y = \{k < \omega : \langle \omega; X[G] \rangle \models \Phi(k)\}$; then $y \in \mathbf{L}[G]$, $y \subseteq \omega$. We claim that in fact y belongs to \mathbf{L} , and hence to $X[G]$.

Let \Vdash be the forcing relation associated with \mathbf{P} . In particular, if $p \in \mathbf{P}$ and ψ is a parameter-free formula then $p \Vdash \psi$ iff ψ holds in any \mathbf{P} -generic extension $\mathbf{L}[H]$ of \mathbf{L} such that $p \in H$.

Let \underline{G} be a canonical \mathbf{P} -name for G . We assert that

$$y = \{k < \omega : \Lambda^\omega \Vdash \langle \omega; X[\underline{G}] \rangle \models \Phi(k)\}. \quad (1)$$

Indeed assume that the condition Λ^ω \mathbf{P} -forces “ $\langle \omega; X[\underline{G}] \rangle \models \Phi(k)$ ”. But $\Lambda^\omega \in G$ since Λ^ω is the weakest condition in \mathbf{P} . Therefore $\langle \omega; X[G] \rangle \models \Phi(k)$ by the forcing theorem, thus $k \in y$, as required.

To prove the converse, assume that $k \in y$. Then by the forcing theorem there is a condition $p \in G$ forcing “ $\langle \omega; X[\underline{G}] \rangle \models \Phi(k)$ ”. We claim that then Λ^ω forces the same as well.

Indeed otherwise there is a condition $q \in \mathbf{P}$ which forces “ $\langle \omega; X[\underline{G}] \rangle \models \neg \Phi(k)$ ”. There is a permutation $\pi \in \mathbf{Perm}$ satisfying $|r| \cap |p| = \emptyset$, where $r = \pi q \in \mathbf{P}$. We claim that r forces “ $\langle \omega; X[\underline{G}] \rangle \models \neg \Phi(k)$ ”. Indeed assume that $H \subseteq \mathbf{P}$ is a set \mathbf{P} -generic over \mathbf{L} , and $r \in H$. We have to prove that $\langle \omega; X[H] \rangle \models \neg \Phi(k)$. The set $K = \{\pi r' : r' \in H\}$ is \mathbf{P} -generic over \mathbf{L} along with H since $\pi \in \mathbf{L}$. Moreover K contains q . It follows that $\langle \omega; X[K] \rangle \models \neg \Phi(k)$ by the forcing theorem and the choice of q . However the sequence $\langle x_i[K] \rangle_{i < \omega}$ is equal to the permutation of the sequence $\langle x_i[H] \rangle_{i < \omega}$ by π . It follows that $X[H] = X[K]$, and hence $\langle \omega; X[H] \rangle \models \neg \Phi(k)$, as required. Thus indeed r forces “ $\langle \omega; X[\underline{G}] \rangle \models \neg \Phi(k)$ ”.

However p forces “ $\langle \omega; X[\underline{G}] \rangle \models \Phi(k)$ ”, and p, r are compatible in \mathbf{P} because $|r| \cap |p| = \emptyset$. This is a contradiction.

We conclude that Λ^ω forces $\langle \omega; X[\underline{G}] \rangle \models \Phi(k)$, and this completes the proof of (1).

But it is known that the forcing relation \Vdash is expressible in \mathbf{L} , the ground model. Therefore it follows from (1) that $y \in \mathbf{L}$, hence $y \in X[G]$, as required. \square

4 Generalized Sacks iterations

Here we begin the proof of Theorem 1.2. The proof involves the engine of generalized product/iterated Sacks forcing developed in [7, 8] on the base of earlier

papers [2, 5] and others. We still consider the constructible universe \mathbf{L} as the ground model for the extension, and define, in \mathbf{L} , the set

$$\mathbf{I} = (\omega_1 \times 2^{<\omega}) \cup \omega_1; \quad \mathbf{I} \in \mathbf{L}, \quad (2)$$

partially ordered so that $\langle \gamma, s \rangle \preceq \langle \beta, t \rangle$ iff $\gamma = \beta$ and $s \subseteq t$ in $2^{<\omega}$, while the ordinals in ω_1 (the second part of \mathbf{I}) remain \preceq -incomparable.

Our plan is to define a product/iterated generic Sacks extension $\mathbf{L}[\vec{a}]$ of \mathbf{L} by an array $\vec{a} = \langle a_i \rangle_{i \in \mathbf{I}}$ of reals $a_i \in 2^\omega$, in which the structure of “sacksness” is determined by this set \mathbf{I} , so that in particular each a_i is Sacks-generic over the submodel $\mathbf{L}[\langle a_j \rangle_{j \prec i}]$.

Then we define the set $\mathbf{J} \in \mathbf{L}[\vec{a}]$ of all elements $i \in \mathbf{I}$ such that:

- either $i = \langle \gamma, 0^m \rangle$, where $\gamma < \omega_1$ and $m < \omega$,
- or $i = \langle \gamma, 0^m \wedge 1 \rangle$, where $\gamma < \omega_1$ and $m < \omega$, $a_\gamma(m) = 1$.

This any $i = \langle \gamma, 0^m \rangle \in \mathbf{J}$ is a splitting node in \mathbf{J} iff $a_\gamma(m) = 1$, or in other words

$$a_\gamma(m) = 1 \quad \text{iff} \quad \langle \gamma, 0^m \rangle \text{ is a splitting node in } \mathbf{J}, \quad (3)$$

We’ll finally prove that the according set

$$M = \mathcal{P}(\omega) \cap \bigcup_{i_1, \dots, i_n \in \mathbf{J}} \mathbf{L}[a_{i_1}, \dots, a_{i_n}] \quad (4)$$

leads to the model $\langle \omega; M \rangle$ for Theorem 1.2. The reals a_γ will not belong to M by the choice of \mathbf{J} , but will be definable in $\langle \omega; M \rangle$ (with $a_{\langle \gamma, \Lambda \rangle} \subseteq \omega$ as a parameter) via the characterization of the splitting nodes in \mathbf{J} by (3).

5 Iterated perfect sets

Arguing in \mathbf{L} in this section, we define $\mathbf{I} = \langle \mathbf{I}; \preceq \rangle$ as above.

Let \mathfrak{E} be the set of all countable (including finite) sets $\zeta \subseteq \mathbf{I}$.

If $\zeta \in \mathfrak{E}$ then IS_ζ is the set of all initial segments of ζ .

Greek letters $\xi, \eta, \zeta, \vartheta$ will denote sets in \mathfrak{E} .

Characters i, j are used to denote *elements* of \mathbf{I} .

For any $i \in \zeta \in \mathfrak{E}$, we consider initial segments $\zeta[\prec i] = \{j \in \zeta : j \prec i\}$ and $\zeta[\neq i] = \{j \in \zeta : j \neq i\}$, and $\zeta[\preceq i], \zeta[\neq i]$ defined analogously.

Further, ω^ω is the *Baire space*. Points of ω^ω will be called *reals*.

Let $\mathcal{D} = 2^\omega \subseteq \omega^\omega$ be the *Cantor space*. For any countable set ξ , \mathcal{D}^ξ is the product of ξ -many copies of \mathcal{D} with the product topology. Then every \mathcal{D}^ξ is a compact space, homeomorphic to \mathcal{D} itself unless $\xi = \emptyset$.

Assume that $\eta \subseteq \xi \in \mathfrak{E}$. If $x \in \mathcal{D}^\xi$ then let $x \upharpoonright \eta \in \mathcal{D}^\eta$ denote the usual restriction. If $X \subseteq \mathcal{D}^\xi$ then let $X \upharpoonright \eta = \{x \upharpoonright \eta : x \in X\}$. To save space, let $X \upharpoonright \prec i$ mean $X \upharpoonright \xi[\prec i]$, $X \upharpoonright \neq i$ mean $X \upharpoonright \xi[\neq i]$, etc.

But if $Y \subseteq \mathcal{D}^n$ then we put $Y \upharpoonright^{-1} \xi = \{x \in \mathcal{D}^\xi : x \upharpoonright \eta \in Y\}$.

To describe the idea behind the definition of iterated perfect sets, recall that the Sacks forcing consists of perfect subsets of \mathcal{D} , that is, sets of the form $H''\mathcal{D} = \{H(a) : a \in \mathcal{D}\}$, where $H : \mathcal{D} \xrightarrow{\text{onto}} X$ is a homeomorphism.

To get a product Sacks model, with two factors (the case of a two-element unordered set as the length of iteration), we have to consider sets $X \subseteq \mathcal{D}^2$ of the form $X = H''\mathcal{D}^2$ where H , a homeomorphism defined on \mathcal{D}^2 , splits in obvious way into a pair of one-dimensional homeomorphisms.

To get an iterated Sacks model, with two stages of iteration (the case of a two-element ordered set as the length of iteration), we have to consider sets $X \subseteq \mathcal{D}^2$ of the form $X = H''\mathcal{D}^2$, where H , a homeomorphism defined on \mathcal{D}^2 , satisfies the following: if $H(a_1, a_2) = \langle x_1, x_2 \rangle$ and $H(a'_1, a'_2) = \langle x'_1, x'_2 \rangle$ then $a_1 = a'_1 \iff x_1 = x'_1$.

The combined product/iteration case results in the following definition.

Definition 5.1 (iterated perfect sets, [7, 8]). For any $\zeta \in \mathfrak{E}$, \mathbf{Perf}_ζ is the collection of all sets $X \subseteq \mathcal{D}^\zeta$ such that there is a homeomorphism $H : \mathcal{D}^\zeta \xrightarrow{\text{onto}} X$ satisfying

$$x_0 \upharpoonright \xi = x_1 \upharpoonright \xi \iff H(x_0) \upharpoonright \xi = H(x_1) \upharpoonright \xi$$

for all $x_0, x_1 \in \text{dom } H$ and $\xi \in \mathfrak{E}$, $\xi \subseteq \zeta$. Homeomorphisms H satisfying this requirement will be called *projection-keeping*. In other words, sets in \mathbf{Perf}_ζ are images of \mathcal{D}^ζ via projection-keeping homeomorphisms. \square

Remark 5.2. Note that \emptyset , the empty set, formally belongs to \mathfrak{E} , and then $\mathcal{D}^\emptyset = \{\emptyset\}$, and we easily see that $\mathbb{1} = \{\emptyset\}$ is the only set in \mathbf{Perf}_\emptyset . \square

For the convenience of the reader, we now present five lemmas on sets in \mathbf{Perf}_ζ established in [7, 8].

Lemma 5.3 (Proposition 4 in [7]). *Let $\zeta \in \mathfrak{E}$. Every set $X \in \mathbf{Perf}_\zeta$ is closed and satisfies the following properties:*

- P-1. *If $\mathbf{i} \in \zeta$ and $z \in X \upharpoonright_{<\mathbf{i}}$ then $D_{Xz}(\mathbf{i}) = \{x(\mathbf{i}) : x \in X \wedge x \upharpoonright_{<\mathbf{i}} = z\}$ is a perfect set in \mathcal{D} .*
- P-2. *If $\xi \in \text{IS}_\zeta$, and a set $X' \subseteq X$ is open in X (in the relative topology) then the projection $X' \upharpoonright \xi$ is open in $X \upharpoonright \xi$. In other words, the projection from X to $X \upharpoonright \xi$ is an open map.*
- P-3. *If $\xi, \eta \in \text{IS}_\zeta$, $x \in X \upharpoonright \xi$, $y \in X \upharpoonright \eta$, and $x \upharpoonright (\xi \cap \eta) = y \upharpoonright (\xi \cap \eta)$, then $x \cup y \in X \upharpoonright (\xi \cup \eta)$.*

Proof (sketch). Clearly \mathcal{D}^ζ satisfies P-1, P-2, P-3, and one easily shows that projection-keeping homeomorphisms preserve the requirements. \square

Lemma 5.4 (Lemma 6 in [7]). *If $\zeta \in \mathfrak{E}$, $X \in \mathbf{Perf}_\zeta$, $\xi \in \text{IS}_\zeta$, then $X \upharpoonright \xi \in \mathbf{Perf}_\xi$.*

Lemma 5.5 (Lemma 8 in [7]). *If $\zeta \in \Xi$, $X \in \mathbf{Perf}_\zeta$, a set $X' \subseteq X$ is open in X , and $x_0 \in X'$, then there is a set $X'' \in \mathbf{Perf}_\zeta$, $X'' \subseteq X'$, clopen in X and containing x_0 .*

Lemma 5.6 (Lemma 10 in [7]). *Suppose that $\zeta \in \Xi$, $\eta \in \mathbf{IS}_\zeta$, $X \in \mathbf{Perf}_\zeta$, $Y \in \mathbf{Perf}_\eta$, and $Y \subseteq X \upharpoonright \eta$. Then $Z = X \cap (Y \upharpoonright^{-1} \zeta)$ belongs to \mathbf{Perf}_ζ .*

Lemma 5.7 (Lemma 10 in [8]). *Suppose that $\zeta \in \Xi$, $\xi \subseteq \zeta$, $X \in \mathbf{Perf}_\xi$. Then $X \upharpoonright^{-1} \zeta$ belongs to \mathbf{Perf}_ζ .*

6 The forcing and the basic extension

This section introduces the forcing notion we consider and the according generic extension called the basic extension.

We continue to argue in \mathbf{L} . Recall that a partially ordered set $I \in \mathbf{L}$ is defined by (2) in Section 4, and Ξ is the set of all at most countable initial segments $\xi \subseteq I$ in \mathbf{L} . For any $\zeta \in \Xi$, let $\mathbb{P}_\zeta = (\mathbf{Perf}_\zeta)^\mathbf{L}$.

The set $\mathbb{P} = \mathbb{P}_I = \bigcup_{\zeta \in \Xi} \mathbb{P}_\zeta \in \mathbf{L}$ will be the *forcing notion*.

To define the order, we put $\|X\| = \zeta$ whenever $X \in \mathbb{P}_\zeta$. Now we set $X \leq Y$ (i.e. X is *stronger* than Y) iff $\zeta = \|Y\| \subseteq \|X\|$ and $X \upharpoonright \zeta \subseteq Y$.

Remark 6.1. We may note that the set $\mathbb{1} = \{\emptyset\}$ as in Remark 5.2 belongs to \mathbb{P} and is the \leq -largest (i.e., the weakest) element of \mathbb{P} . \square

Now let $G \subseteq \mathbb{P}$ be a \mathbb{P} -generic set (filter) over \mathbf{L} .

Remark 6.2. If $X \in \mathbb{P}_\zeta$ in \mathbf{L} then X is not even a closed set in \mathcal{D}^ζ in $\mathbf{L}[G]$. However we can transform it to a perfect set in $\mathbf{L}[G]$ by the closure operation. Indeed the topological closure $X^\#$ of such a set X in \mathcal{D}^ζ taken in $\mathbf{L}[G]$ belongs to \mathbf{Perf}_ζ from the point of view of $\mathbf{L}[G]$. \square

It easily follows from Lemma 5.5 that there exists a unique array $\mathbf{a}[G] = \langle \mathbf{a}_i[G] \rangle_{i \in I}$, all $\mathbf{a}_i[G]$ being elements of 2^ω , such that $\mathbf{a}[G] \upharpoonright \xi \in X^\#$ whenever $X \in G$ and $\|X\| = \xi \in \Xi$. Then $\mathbf{L}[G] = \mathbf{L}[\langle \mathbf{a}_i[G] \rangle_{i \in I}] = \mathbf{L}[\mathbf{a}[G]]$ is a \mathbb{P} -generic extension of \mathbf{L} .

Theorem 6.3 (Theorems 24, 31 in [7]). *Every cardinal in \mathbf{L} remains a cardinal in $\mathbf{L}[G]$. Every $\mathbf{a}_i[G]$ is Sacks generic over the model $\mathbf{L}[\mathbf{a}[G] \upharpoonright \prec_i]$.*

We now present several lemmas on reals in \mathbb{P} -generic models $\mathbf{L}[G]$, established in [7]. In the lemmas, we let $G \subseteq \mathbb{P}$ be a set \mathbb{P} -generic over \mathbf{L} .

Lemma 6.4 (Lemma 22 in [7]). *Suppose that sets $\eta, \xi \in \Xi$ satisfy $\forall j \in \eta \exists i \in \xi (j \preceq i)$. Then $\mathbf{a}[G] \upharpoonright \eta \in \mathbf{L}[\mathbf{a}[G] \upharpoonright \xi]$.*

Lemma 6.5 (Lemma 26 in [7]). *Suppose that $\mathbf{K} \in \mathbf{L}$ is an initial segment in I , and $i \in I \setminus \mathbf{K}$. Then $\mathbf{a}_i[G] \notin \mathbf{L}[\mathbf{a}[G] \upharpoonright \mathbf{K}]$.*

Lemma 6.6 (Corollary 27 in [7]). *If $i \neq j$ then $\mathbf{a}_i[G] \neq \mathbf{a}_j[G]$ and even $\mathbf{L}[\mathbf{a}_i[G]] \neq \mathbf{L}[\mathbf{a}_j[G]]$.* \square

Lemma 6.7 (Lemma 29 in [7]). *If $K \in \mathbf{L}$ is an initial segment of I , and r is a real in $\mathbf{L}[G]$, then either $r \in \mathbf{L}[x \upharpoonright K]$ or there is $i \notin K$ such that $\mathbf{a}_i[G] \in \mathbf{L}[r]$.*

We apply the lemmas in the proof of the next theorem. Let $\leq_{\mathbf{L}}$ denote the Gödel wellordering on 2^ω , so that $x \leq_{\mathbf{L}} y$ iff $x \in \mathbf{L}[y]$. Let $x <_{\mathbf{L}} y$ mean that $x \leq_{\mathbf{L}} y$ but $y \not\leq_{\mathbf{L}} x$, and $x \equiv_{\mathbf{L}} y$ mean that $x \leq_{\mathbf{L}} y$ and $y \leq_{\mathbf{L}} x$.

Theorem 6.8. *Assume that $i \in I$ and $r \in \mathbf{L}[G] \cap 2^\omega$. Then*

- (i) *if $j \in I$ and $j \preceq i$ then $\mathbf{a}_j[G] \leq_{\mathbf{L}} \mathbf{a}_i[G]$;*
- (ii) *if $j \in I$ and $j \not\preceq i$ then $\mathbf{a}_j[G] \not\leq_{\mathbf{L}} \mathbf{a}_i[G]$;*
- (iii) *if $r \leq_{\mathbf{L}} \mathbf{a}_i[G]$ then $r \in \mathbf{L}$ or $r \equiv_{\mathbf{L}} \mathbf{a}_j[G]$ for some $j \in I$, $j \preceq i$;*
- (iv) *if $i = \langle \gamma, s \rangle \in I$, $e = 0, 1$, and $i \frown e = \langle \gamma, s \frown e \rangle$ then $\mathbf{a}_{i \frown e}[G]$ is a **true successor** of $\mathbf{a}_i[G]$ in the sense that $\mathbf{a}_i[G] <_{\mathbf{L}} \mathbf{a}_{i \frown e}[G]$ and any real $y \in 2^\omega$ satisfies $y <_{\mathbf{L}} \mathbf{a}_{i \frown e}[G] \implies y \leq_{\mathbf{L}} \mathbf{a}_i[G]$;*
- (v) *if $i = \langle \gamma, s \rangle \in I$, and $x \in 2^\omega \cap \mathbf{L}[G]$ is a true successor of $\mathbf{a}_i[G]$ in the sense of (iv), then there is $e = 0$ or 1 such that $x \equiv_{\mathbf{L}} \mathbf{a}_{i \frown e}[G]$.*

Proof. (i) Apply Lemma 6.4 with $\eta = \{j\}$ and $\xi = \{i\}$.

(ii) Apply Lemma 6.5 with $K = [\preceq i]$.

(iii) If there are elements $j \in \mathcal{I}$, $j \preceq i$, such that $\mathbf{a}_j[G] \in \mathbf{L}[r]$, then let j be the largest such one, and let $\xi = [\preceq j]$ (a finite initial segment of I). Then, by Lemma 6.7, either $r \in \mathbf{L}[\mathbf{a}[G] \upharpoonright \xi]$, or there is $i' \notin \xi$ such that $\mathbf{a}_{i'}[G] \in \mathbf{L}[r]$.

In the “either” case, we have $r \in \mathbf{L}[\mathbf{a}_j[G]]$ by (i), so that $\mathbf{L}[r] = \mathbf{L}[\mathbf{a}_j[G]]$ by the choice of j . In the “or” case we have $\mathbf{a}_{i'}[G] \in \mathbf{L}[\mathbf{a}_i[G]]$, hence $i' \preceq i$ by (ii). But this contradicts the choice of j and i' .

Finally if there is no $j \in \mathcal{I}$, $j \preceq i$, such that $\mathbf{a}_j[G] \in \mathbf{L}[r]$, then the same argument with $\xi = \emptyset$ gives $r \in \mathbf{L}$.

(iv) The relation $\mathbf{a}_j[G] <_{\mathbf{L}} \mathbf{a}_{i \frown e}[G]$ is implied by Lemmas 6.4 and 6.5. If now $y <_{\mathbf{L}} \mathbf{a}_{i \frown e}[G]$ then $y \in \mathbf{L}$ or $y \equiv_{\mathbf{L}} \mathbf{a}_j[G]$ for some $j \preceq i \frown e$ by (iii), and in the latter case in fact $j \prec i \frown e$, hence $j \preceq i$, and then $y \leq_{\mathbf{L}} \mathbf{a}_i[G]$.

(v) By (iv), it suffices to prove that $x \leq_{\mathbf{L}} \mathbf{a}_{i \frown 0}[G]$ or $x \leq_{\mathbf{L}} \mathbf{a}_{i \frown 1}[G]$. Assume that $x \not\leq_{\mathbf{L}} \mathbf{a}_{i \frown 0}[G]$. Then by Lemma 6.7 there is an element $j \in I$ such that $j \not\preceq i \frown 0$ and $\mathbf{a}_j[G] \leq_{\mathbf{L}} x$. If $\mathbf{a}_j[G] <_{\mathbf{L}} x$ strictly then $\mathbf{a}_j[G] \leq_{\mathbf{L}} \mathbf{a}_i[G]$ by the true successor property, hence $i_0 \preceq i$, contrary to $i_0 \not\preceq i \frown 0$, see above. Therefore in fact $\mathbf{a}_{i_0}[G] \equiv_{\mathbf{L}} x$. Then we must have $i_0 = i \frown 0$ or $i_0 = i \frown 1$ as x is a true successor, but then $i_0 = i \frown 1$, as $x \not\leq_{\mathbf{L}} \mathbf{a}_{i \frown 0}[G]$ was assumed, and we are done. \square

7 The subextension

Following the arguments above, assume that $G \subseteq \mathbb{P}$ is a set \mathbb{P} -generic over \mathbf{L} , and consider the set $\mathbf{J}[G] \in \mathbf{L}[G]$ of all elements $\mathbf{i} \in \mathbf{I}$ such that either $\mathbf{i} = \langle \gamma, 0^m \rangle$, where $\gamma < \omega_1$ and $m < \omega$, or $\mathbf{i} = \langle \gamma, 0^m \hat{\ } 1 \rangle$, where $\gamma < \omega_1$ and $m < \omega$, $\mathbf{a}_\gamma[G](m) = 1$. Following (4), we define

$$M[G] = \mathcal{P}(\omega) \cap \bigcup_{\mathbf{i}_1, \dots, \mathbf{i}_n \in \mathbf{J}[G]} \mathbf{L}[a_{\mathbf{i}_1}[G], \dots, a_{\mathbf{i}_n}[G]], \quad (5)$$

Lemma 7.1. *If $\mathbf{i} \notin \mathbf{J}[G]$ then $\mathbf{a}_\mathbf{i}[G] \notin M[G]$.*

Proof. This is not immediately a case of Lemma 6.5 because $\mathbf{J}[G] \notin \mathbf{L}$. However the set $\mathbf{K} = \{\mathbf{j} \in \mathbf{I} : \mathbf{i} \not\leq \mathbf{j}\}$ belongs to \mathbf{L} and satisfies $\mathbf{J}[G] \subseteq \mathbf{K} \subseteq \mathbf{I}$. We have $\mathbf{i} \notin \mathbf{K}$, and hence $\mathbf{a}_\mathbf{i}[G] \notin \mathbf{L}[\mathbf{a}[G] \upharpoonright \mathbf{K}]$ by Lemma 6.5. On the other hand, we easily check $X \subseteq \mathbf{L}[\mathbf{a}[G] \upharpoonright \mathbf{K}]$, and we are done. \square

We are going to prove that $\langle \omega; M[G] \rangle$ is a model of $\mathbf{PA}_2^* + \mathbf{CA}(\Sigma_2^1)$, but the full \mathbf{CA} fails in $\langle \omega; M[G] \rangle$.

Part 1: $\langle \omega; M[G] \rangle$ is a model of all axioms of \mathbf{PA}_2 except for \mathbf{CA} , trivial.

Part 2: $\langle \omega; M[G] \rangle$ is a model of $\mathbf{CA}(\Sigma_2^1)$ (with parameters). This is also easy by the Shoenfield absoluteness theorem.

Part 3: $\langle \omega; M[G] \rangle$ fails to satisfy the full \mathbf{CA} . Here we need some work. Let $\gamma < \omega_1^{\mathbf{L}}$, so that both γ and each pair $\langle \gamma, s \rangle$, $s \in 2^{<\omega}$, belong to \mathbf{I} by (2) in Section 4, in particular $\mathbf{i}_0 = \langle \gamma, \Lambda \rangle \in \mathbf{I}$, where Λ is the empty tuple. In addition γ (as an element of \mathbf{I}) does not belong to $\mathbf{J}[G]$. Our plan is to prove that $\mathbf{a}_\gamma[G] \notin M[G]$ but $\mathbf{a}_\gamma[G]$ is definable in $\langle \omega; M[G] \rangle$.

Subpart 3.1: $\mathbf{a}_\gamma[G] \notin M[G]$ by Lemma 7.1 just because $\gamma \notin \mathbf{J}[G]$.

Subpart 3.2: $\mathbf{a}_\gamma[G]$ is definable in $\langle \omega; M[G] \rangle$ with $\mathbf{a}_{\mathbf{i}_0}[G]$ as a parameter, where $\mathbf{i}_0 = \langle \gamma, \Lambda \rangle \in \mathbf{J}[G]$. Namely we claim that for any $m < \omega$:

$$\mathbf{a}_\gamma[G](m) = 1 \quad \text{iff} \quad \text{there is an array of reals } b_0, b_1, \dots, b_m, b_{m+1} \text{ and } b'_{m+1} \text{ in } 2^\omega \text{ such that } b_0 = \mathbf{a}_{\mathbf{i}_0}, \text{ each } b_{k+1} \text{ is a true successor of } b_k \text{ (} k \leq m \text{), } b'_{m+1} \text{ is a true successor of } b_m \text{ as well, and } b'_{m+1} \neq_{\mathbf{L}} b_{m+1}. \quad (6)$$

The formula in the right-hand side of (6) is based on the Gödel canonical Σ_2^1 formula for $\leq_{\mathbf{L}}$, which is absolute for $M[G]$ by the definition of $M[G]$. Therefore (6) implies that $\mathbf{a}_\gamma[G]$ is definable in $\langle \omega; M[G] \rangle$ with $\mathbf{a}_{\mathbf{i}_0}[G]$ as a parameter. Thus it remains to establish (6).

Direction \implies . Assume that $\mathbf{a}_\gamma[G](m) = 1$. Then $\mathbf{J}[G]$ contains the elements $\mathbf{i}_k = \langle \gamma, 0^k \rangle$, $k \leq m+1$, along with an element $\mathbf{i}'_{m+1} = \langle \gamma, 0^m \hat{\ } 1 \rangle$. Therefore the reals $b_k = \mathbf{a}_{\mathbf{i}_k}[G]$, $k \leq m+1$, and $b'_{m+1} = \mathbf{a}_{\mathbf{i}'_{m+1}}[G]$ belong to $M[G]$. Now

Theorem 6.8(iv),(ii) implies that the reals b_k and b'_{m+1} satisfy the right-hand side of (6), as required.

Direction \Leftarrow . Assume that the reals b_k , $k \leq m+1$, and b'_{m+1} satisfy the right-hand side of (6). By Theorem 6.8(v), there is an array of bits e_1, \dots, e_m, e_{m+1} and e'_{m+1} such that $b_k = \mathbf{a}_{i_k}[G]$ for all $k \leq m+1$ and $b'_{m+1} = \mathbf{a}_{i'_{m+1}}[G]$, where $i_k = \langle \gamma, \langle e_1, \dots, e_k \rangle \rangle$ and $i'_{m+1} = \langle \gamma, \langle e_1, \dots, e_m, e'_{m+1} \rangle \rangle$.

However we must have $i_k \in \mathbf{J}[G]$ for all $k \leq m+1$, and $i'_{m+1} \in \mathbf{J}[G]$, by Lemma 7.1, since the reals b_k and b'_{m+1} belong to $M[G]$. Then obviously $e_1 = \dots = e_m = 0$ while $e_{m+1} = 0$ and $e'_{m+1} = 1$ or vice versa $e_{m+1} = 1$ and $e'_{m+1} = 0$. In other words, the elements $\langle \gamma, 0^{m+1} \rangle$ and $\langle \gamma, 0^m \wedge 1 \rangle$ belong to $\mathbf{J}[G]$. This implies $\mathbf{a}_\gamma[G](m) = 1$.

Part 4: $\langle \omega; M[G] \rangle$ satisfies the parameter-free schema \mathbf{CA}^* . This is rather similar to the verification of \mathbf{CA}^* in $\langle \omega; X[G] \rangle$ in Section 3.

Assume that $\Phi(k)$ is a parameter-free $\mathcal{L}(\mathbf{PA}_2)$ formula with k the only free variable. Consider the set $y = \{k < \omega : \langle \omega; M[G] \rangle \models \Phi(k)\}$; then $y \in \mathbf{L}[G]$, $y \subseteq \omega$. We claim that y even belongs to \mathbf{L} , and hence to $M[G]$.

Let \Vdash be the forcing relation associated with \mathbb{P} , over \mathbf{L} as the ground model. Thus if $X \in \mathbb{P}$ and $k < \omega$ then $X \Vdash \Phi(k)$ iff $\Phi(k)$ holds in any \mathbb{P} -generic extension $\mathbf{L}[H]$ of \mathbf{L} such that $X \in H$.⁴ Let \underline{G} be a canonical \mathbb{P} -name for G . We assert that

$$y = \{k < \omega : \mathbb{1} \Vdash \langle \omega; M[\underline{G}] \rangle \models \Phi(k)\}. \quad (7)$$

(See Remark 6.1 on $\mathbb{1}$.)

In the nontrivial direction, assume that $k \in y$. Then by the forcing theorem there is a condition $X \in G$ forcing $\langle \omega; M[\underline{G}] \rangle \models \Phi(k)$. We claim that then $\mathbb{1}$ forces the same as well.

To prove this reduction, we define, **still in \mathbf{L}** , the set $\mathbf{Perm} \in \mathbf{L}$ that consists of all bijections $\pi : \omega_1 \xrightarrow{\text{onto}} \omega_1$ such that $\pi = \pi^{-1}$ and *the domain of nontriviality* $|\pi| = \{\alpha : \pi(\alpha) \neq \alpha\}$ is at most countable, *i.e.*, bounded in ω_1 . Any $\pi \in \mathbf{Perm}$ acts on:

- elements $i = \gamma$ or $i = \langle \gamma, s \rangle$ of \mathbf{I} , by $\pi i = \pi(\gamma)$, resp. $i = \langle \pi(\gamma), s \rangle$;
- maps g with $\text{dom } g \subseteq \mathbf{I}$, by $\text{dom}(\pi g) = \pi'' \text{dom } g$ and $(\pi g)(\pi(\alpha)) = g(\alpha)$ for all $\alpha \in \text{dom } g$;
- thus if $\xi \subseteq \mathbf{I}$ and $x \in \mathcal{D}^\xi$ then $\pi x \in \mathcal{D}^{\pi''\xi}$ and $(\pi x)(\pi(\alpha)) = x(\alpha)$;
- sets $X \in \mathbf{Perf}_\xi$, $\xi \in \mathbf{\Xi}$, by $\pi X = \{\pi x : x \in X\} \in \mathbf{Perf}_{\pi''\xi}$.

We return to the nontrivial direction \implies of (7), where we have to prove that the condition $\mathbb{1}$ forces “ $\langle \omega; M[\underline{G}] \rangle \models \Phi(k)$ ”. Let this be not the case.

⁴ See Kunen [10] on forcing, especially Section IV.6 there on the “forcing over the universe” approach.

Then there is a condition $Y \in \mathbb{P}$ which forces “ $\langle \omega; M[\underline{G}] \rangle \models \neg \Phi(k)$ ”. There is a permutation $\pi \in \mathbf{Perm}$ satisfying $\|Z\| \cap \|X\| = \emptyset$, where $Z = \pi Y \in \mathbb{P}$. We claim that Z forces “ $\langle \omega; M[\underline{G}] \rangle \models \neg \Phi(k)$ ”. Indeed assume that $H \subseteq \mathbb{P}$ is a set \mathbb{P} -generic over \mathbf{L} , and $Z \in H$. We have to prove that $\langle \omega; M[H] \rangle \models \neg \Phi(k)$. The set $K = \{\pi Z' : Z' \in H\}$ is \mathbf{P} -generic over \mathbf{L} along with H since $\pi \in \mathbf{L}$. Moreover K contains Y . It follows that $\langle \omega; M[K] \rangle \models \neg \Phi(k)$ by the forcing theorem and the choice of Y .

However the array $\mathbf{a}[K]$ is equal to the permutation of the array $\mathbf{a}[H]$ by π . It follows that $M[H] = M[K]$, and hence $\langle \omega; M[H] \rangle \models \neg \Phi(k)$, as required. Thus indeed Z forces “ $\langle \omega; M[\underline{G}] \rangle \models \neg \Phi(k)$ ”.

Recall that X forces “ $\langle \omega; M[\underline{G}] \rangle \models \Phi(k)$ ”. On the other hand, X, Z are compatible in \mathbb{P} because $\|Z\| \cap \|X\| = \emptyset$. This is a contradiction.

We conclude that $\mathbf{1}$ forces “ $\langle \omega; M[\underline{G}] \rangle \models \Phi(k)$ ”, and this completes the proof of (7). But it is known that the forcing relation \Vdash is expressible in \mathbf{L} , the ground model. Therefore it follows from (7) that $y \in \mathbf{L}$, hence $y \in M[G]$, as required.

8 Remarks and questions

Here we present three questions related to possible extensions of Theorem 1.2.

Problem 8.1. Is the parameter-free countable choice schema \mathbf{AC}^* in the language $\mathcal{L}(\mathbf{PA}_2)$ true in the models $\langle \omega; M[G] \rangle$ defined in Section 7?

Problem 8.2. Can we sharpen the result of Theorem 1.2 by specifying that $\mathbf{CA}(\Sigma_3^1)$ is violated? The combination $\mathbf{CA}(\Sigma_2^1)$ plus $\neg \mathbf{CA}(\Sigma_3^1)$ would be optimal. The counterexample to \mathbf{CA} defined in Section 7 (Part 3) definitely is more complex than Σ_3^1 .

Problem 8.3. As a generalization of the above, prove that, for any $n \geq 2$, $\mathbf{PA}_2^* + \mathbf{CA}(\Sigma_n^1)$ does not imply $\mathbf{CA}(\Sigma_{n+1}^1)$. In this case, we'll be able to conclude that the full schema \mathbf{CA} is not finitely axiomatizable over \mathbf{PA}_2^* . Compare to Problem 9 in [1, § 11].

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