

## BOREL OD SETS OF REALS ARE OD-BOREL IN SOME SIMPLE MODELS

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**ABSTRACT.** It is true in the Cohen, Solovay-random, and Sacks generic extensions that every ordinal-definable Borel set of reals has a Borel code in the ground model, and hence if non-empty, then has an element in the ground model.

### 1. INTRODUCTION

It is known from [9] that for each lightface  $\Delta_1^1$  set  $X$ , its Borel class is witnessed by a lightface  $\Delta_1^1$  code. This *effective Borel coding* property is not necessarily true for more general definability classes instead of  $\Delta_1^1$ . For instance there are models of **ZFC** in which there exists a countable, hence  $\mathbf{F}_\sigma$ , lightface  $\Pi_2^1$  non-empty set  $X$  of reals with no OD<sup>1</sup> elements [3, 4], and such a set  $X$  definitely has no OD  $\mathbf{F}_\sigma$  code. These models make use of very non-homogeneous forcing notions.<sup>2</sup> Therefore one may expect that homogeneous forcing notions generally yield opposite results. Working in this direction, we prove here the following theorem.

**Theorem 1.** *Let  $a$  be either (A) a Cohen-generic real or (B) a Solovay-random real or (C) a Sacks real over the set universe  $\mathbf{V}$ . Then it is true in  $\mathbf{V}[a]$  that if  $X \subseteq 2^\omega$  is a Borel OD set, then it has a Borel code in  $\mathbf{V}$  of the same ordinal level.*

One may expect such a theorem to be true for other suitably homogeneous generic models like e.g. the dominating forcing extensions. However it does not seem to be an easy task to manufacture a proof of sufficient generality because of various ad hoc arguments in the proofs below, lacking a common denominator.

### 2. BOREL CODING

We fix any reasonable system of Borel coding, which involves a  $\Pi_1^1$  set  $\mathbf{BC} \subseteq \omega^\omega$  of *Borel codes* and an assignment of a Borel set  $\mathbf{B}_c \subseteq \omega^\omega$  for each  $c \in \mathbf{BC}$ , as e.g. in [6, 2.9] or [10, 3H]. This also includes a pair of  $\Pi_1^1$  sets  $S, S' \subseteq \omega^\omega \times \omega^\omega$  such that we have  $x \in \mathbf{B}_c \iff \langle c, x \rangle \in S \iff \langle c, x \rangle \notin S'$  whenever  $c \in \mathbf{BC}$  and  $x \in \omega^\omega$  are

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<sup>1</sup>OD means *ordinal-definable*, that is, a set which can be defined by an  $\in$ -formula with ordinals as parameters. See Jech [1, Section 13] on this notion.

<sup>2</sup>The model in [4] involves the countable product of Jensen's minimal  $\Delta_3^1$  real forcing [2]. The model in [3] involves a shift-invariant version of Jensen's forcing, and it contains a countable  $\Pi_2^1$  set  $X \subseteq \omega^\omega$  of reals with no OD elements, and  $X$  is an  $\mathbf{E}_0$ -class.

arbitrary. If  $1 \leq \xi < \omega_1$ , then there is a Borel subset  $\mathbf{BC}_\xi \subseteq \mathbf{BC}$  which canonically produces  $\mathbf{\Pi}_\xi^0$  sets, so that  $\{\mathbf{B}_c : c \in \mathbf{BC}_\xi\}$  is equal to the set of all boldface  $\mathbf{\Pi}_\xi^0$  sets  $X \subseteq \omega^\omega$ .

To accordingly code *Borel maps*  $F : \omega^\omega \rightarrow \omega^\omega$ , we let  $\mathbf{FC}$  be the lightface  $\mathbf{\Pi}_1^1$  set of all reals  $h \in \omega^\omega$  such that  $(h)_n \in \mathbf{BC}$ ,  $\forall n$ , where  $(h)_n \in \omega^\omega$  is defined by  $(h)_n(k) = h(2^n(2k + 1) - 1)$  for all  $k$ . If  $h \in \mathbf{FC}$ , then a Borel map  $\vartheta_h : \omega^\omega \rightarrow \omega^\omega$  (a total map with the full domain  $\omega^\omega$ ) is defined so that  $\vartheta_h(x)(n) = k$  iff either  $k \geq 1$  and  $x \in \mathbf{B}_{(h)_n(k)} \setminus \bigcup_{1 \leq \ell < k} \mathbf{B}_{(h)_n(\ell)}$  or  $k = 0$  and  $x \notin \bigcup_{1 \leq \ell} \mathbf{B}_{(h)_n(\ell)}$ .

*Remark 2.* There is another system of Borel codes of the form  $c = \langle T_c, f_c \rangle$ , where  $T_c$  is a well-founded tree and  $f_c$  maps terminal nodes of  $T$  into *Baire intervals* in  $\omega^\omega$ ; see e.g. [12]. If one assumes that  $T_c$  is a tree in  $\omega^{<\omega}$ , then this is fully equivalent to the above system of coding by  $\mathbf{B}_c, c \in \mathbf{BC}$ .

But assuming that  $T_c \subseteq \lambda^{<\omega}$ ,  $\lambda < \omega_1$  leads to new insights, and then, as essentially proved in [12], our Theorem 1 is true also in the Solovay model (the Levy-collapse extension of  $\mathbf{L}$ ) in such a way that the codes  $c = \langle T_c, f_c \rangle$  which witness the Borel class of Borel OD sets belong to  $\mathbf{L}$ , but the trees  $T_c$  may not be countable in  $\mathbf{L}$ .

As for the coding system by  $\mathbf{B}_c, c \in \mathbf{BC} \subseteq \omega^\omega$ , Theorem 1 obviously fails in the Solovay model, the countable set  $X = \omega^\omega \cap \mathbf{L}$  being a counterexample.

### 3. COHEN-GENERIC REALS, CASE A OF THEOREM 1

Let  $\mathbf{Coh} = 2^{<\omega}$  be the Cohen forcing.

*Proof of Theorem 1, case A.* Let  $a_0 \in 2^\omega$  be a real  $\mathbf{Coh}$ -generic over the background set universe  $\mathbf{V}$ . Assume that  $1 \leq \xi < \omega_1 (= \omega_1^{\mathbf{V}})$ , and know that in  $\mathbf{V}[a_0]$  it is true that  $X = \{x \in 2^\omega : \varphi(x)\} \subseteq 2^\omega$  is a  $\mathbf{\Pi}_\xi^0$  set definable by a formula  $\varphi$  with sets in  $\mathbf{V}$  as parameters; this includes the OD case. Suppose to the contrary that there is no Borel code  $c \in \mathbf{BC}_\xi$  in  $\mathbf{V}$  satisfying  $X = \mathbf{B}_c$ .

As  $X$  is  $\mathbf{\Pi}_\xi^0$ , there is a code  $d \in \mathbf{BC}_\xi$  in  $\mathbf{V}[a_0]$  satisfying  $X = \mathbf{B}_d$ . Cohen extensions are known to satisfy the property of *Borel reading of names*; hence  $d = \vartheta_h(a_0)$ , where  $h \in \mathbf{FC}$  in  $\mathbf{V}$ . Thus  $X = \mathbf{B}_{\vartheta_h(a_0)}$ . It follows that there is a condition  $u \in \mathbf{Coh}$  which forces  $\vartheta_h(\dot{a}) \in \mathbf{BC}_\xi$  and  $\{x \in 2^\omega : \varphi(x)\} = \mathbf{B}_{\vartheta_h(a_0)}$ , where  $\dot{a}$  is a canonical name for the Cohen generic real in  $2^\omega$ , and also forces that there is no Borel code  $c \in \mathbf{BC}_\xi$  in  $\mathbf{V}$  satisfying  $\{x \in 2^\omega : \varphi(x)\} = \mathbf{B}_c$ .

**Argue in the universe  $\mathbf{V}$ .** The set  $I_u = \{x \in 2^\omega : u \subset x\}$  is a *Cantor interval*, clopen in  $2^\omega$ . The set  $\{x \in I_u : \vartheta_h(x) \in \mathbf{BC}_\xi\}$  is comeager in  $I_u$  by the choice of  $u$ . It follows that there is a dense  $\mathbf{G}_\delta$  set  $D \subseteq I_u$  such that  $\vartheta_h(x) \in \mathbf{BC}_\xi$  for all  $x \in D$ . Consider the Borel set

$$P = \{\langle x, y \rangle : x \in D \wedge y \in \mathbf{B}_{\vartheta_h(x)}\} \subseteq 2^\omega \times \omega^\omega$$

and the  $\mathbf{\Pi}_1^1$  equivalence relation  $x \mathbf{E} x'$  iff  $P_x = P_{x'}$ , on  $D$ , where as usual  $P_x = \{y : P(x, y)\}$ . As a subset of  $I_u \times I_u$ ,  $\mathbf{E}$  has the Baire property, and so do all  $\mathbf{E}$ -equivalence classes. Thus there is a condition  $v \in \mathbf{Coh}$  which satisfies the requirements of one of the two cases below.

*Case 1 (in  $\mathbf{V}$ ).* All  $\mathbf{E}$ -equivalence classes are meager on  $I_v = \{x \in 2^\omega : v \subset x\}$ . Then the  $\mathbf{\Pi}_1^1$ -set  $W = \{\langle x, x' \rangle : x, x' \in I_v \cap D \wedge x \mathbf{E} x'\}$  is meager in  $I_v \times I_v$  by Ulam–Kuratowski. Therefore  $W$  is covered by an  $\mathbf{F}_\sigma$  meager set  $F \subseteq I_v \times I_v$ .

Fix a transitive model  $\mathfrak{M}$  of a sufficiently large fragment of **ZFC** which contains the code  $h$  and codes for  $D, F$  and is an elementary submodel of the universe  $\mathbf{V}$  w.r.t. all analytic formulas.

**Lemma 3.** *There are reals  $a, b \in I_v$ , Cohen generic over  $\mathbf{V}$ , such that  $\mathbf{V}[a] = \mathbf{V}[b]$  and the pair  $\langle a, b \rangle$  is Cohen  $\times$  Cohen generic over  $\mathfrak{M}$ .*

*Proof.* The set  $P = \{\langle x, x +_2 y \rangle : x, y \in I_v\}$  is non-meager, hence so is the projection  $Z = \{z \in 2^\omega : P^z \text{ non-meager}\}$  by Ulam–Kuratowski, where  $P^z = \{x : \langle x, z \rangle \in P\}$  and  $+_2$  is the componentwise addition mod 2. Let, in  $\mathbf{V}$ ,  $z \in Z$  be Cohen generic over  $\mathfrak{M}$ . Then  $P^z$  is non-meager. Pick a real  $a \in P^z$  Cohen over  $\mathbf{V}$ , hence over  $\mathfrak{M}[z]$ . The pair  $\langle a, z \rangle$  belongs to  $P$  and is Cohen over  $\mathfrak{M}$ ; hence  $z$  is Cohen over  $\mathfrak{M}[a]$ . It follows that  $b = z +_2 a$  is Cohen over  $\mathfrak{M}[a]$ ; thus  $\langle a, b \rangle$  is Cohen over  $\mathfrak{M}$ , and  $a, b \in I_v$  by construction. Finally  $b = z +_2 a$  is Cohen over  $\mathbf{V}$  since so is  $a$  while  $z \in \mathbf{V}$ , and clearly  $\mathbf{V}[a] = \mathbf{V}[b]$ .  $\square$  (Lemma)  $\square$

Consider such a pair of reals  $a, b \in I_v$ . Then  $\langle a, b \rangle \notin F$  since  $F$  is a Borel meager set coded in  $\mathfrak{M}$ . It follows that  $\langle a, b \rangle \notin W$ , hence,  $a \not\equiv b$ , meaning that  $\mathbf{B}_{\mathfrak{g}_h(b)} \neq \mathbf{B}_{\mathfrak{g}_h(a)}$ . But on the other hand  $a, b$  are Cohen generic over  $\mathbf{V}$ ,  $\mathbf{V}[b] = \mathbf{V}[a]$ , and we have  $a, b \in I_v$  by construction. It follows that both  $\mathbf{B}_{\mathfrak{g}_h(b)}$  and  $\mathbf{B}_{\mathfrak{g}_h(a)}$  coincide in  $\mathbf{V}[a]$  with one and the same (since  $\varphi$  has only parameters in  $\mathbf{V}$ ) set  $\{x : \varphi(x)\}$ , which contradicts the above. Thus Case 1 is impossible.

*Case 2 (in  $\mathbf{V}$ ).* At least one of the E-equivalence classes is not meager on  $I_v$ . Then, in  $\mathbf{V}$ , there is a condition  $w \in \mathbf{Coh}$  such that  $v \subseteq w$  and comeager-many points in  $I_w$  are E-equivalent. In other words there exists a particular  $\Pi_\xi^0$  set  $A = \mathbf{B}_c$  with a code  $c \in \mathbf{BC}_\xi$  in  $\mathbf{V}$  such that  $\mathbf{B}_{\mathfrak{g}_h(x)} = \mathbf{B}_c$  for comeager-many  $x \in I_w$ . Then  $w$  Cohen-forces over  $\mathbf{V}$  that  $\mathbf{B}_{\mathfrak{g}_h(\dot{a})} = \mathbf{B}_c$ . But this contradicts the contrary assumption in the beginning of the proof, since  $u \subseteq w$ .

$\square$  (Theorem 1, case A)  $\square$

#### 4. SOLOVAY-RANDOM REALS, CASE B OF THEOREM 1

Let  $\lambda$  be the standard probability Lebesgue measure on  $2^\omega$ . The Solovay-random forcing **Rand** consists of all trees  $T \subseteq 2^{<\omega}$  with no endpoints and no isolated branches and such that the set  $[T] = \{x \in 2^\omega : \forall n (x \upharpoonright n \in T)\}$  has positive measure  $\lambda([T]) > 0$ . The forcing **Rand** depends on the ground model, so that “random over a model  $\mathfrak{M}$ ” will mean “(**Rand**  $\cap$   $\mathfrak{M}$ )-generic over  $\mathfrak{M}$ ”.

Unlike the Cohen-generic case, a random pair of reals is **not** a (**Rand**  $\times$  **Rand**)-generic pair. The notion of a random pair is rather related to forcing by closed sets in  $2^\omega \times 2^\omega$  (or trees which generate them or equivalently Borel sets) of positive product measure (non-null). We’ll make use of the following well-known characterization of Solovay-random pairs.

**Proposition 4.** *Let  $\mathfrak{M}$  be a transitive model of a large fragment of **ZFC**, and let  $a, b \in 2^\omega$ . Then the following three assertions are equivalent:*

- (1) *the pair  $\langle a, b \rangle$  is a random pair over  $\mathfrak{M}$ ;*
- (2)  *$a$  is random over  $\mathfrak{M}$ , and  $b$  is random over  $\mathfrak{M}[a]$ ;*
- (3)  *$b$  is random over  $\mathfrak{M}$ , and  $a$  is random over  $\mathfrak{M}[b]$ .*
- (4)  *$\langle a, b \rangle$  avoids any  $(\lambda \times \lambda)$ -null Borel set  $Q \subseteq 2^\omega \times 2^\omega$  coded in  $\mathfrak{M}$ .*

*Sketch of proof.* Regarding the equivalence  $1 \iff 4$ , see e.g. [1, Lemma 26.4] or V.4.19, V.4.20 in [8], where the 1-dimensional version of the equivalence is established, saying that  $a \in 2^\omega$  is random over  $\mathfrak{M}$  iff  $a$  avoids any  $\lambda$ -null Borel set  $Q \subseteq 2^\omega$  coded in  $\mathfrak{M}$ .

To prove that (1) implies (2), suppose that (2) fails.

If  $a$  is not random over  $\mathfrak{M}$ , then by the same Lemma 26.4 in [1],  $a$  belongs to a null Borel set  $X$  coded in  $\mathfrak{M}$ . Then  $\langle a, b \rangle$  belongs to the  $(\lambda \times \lambda)$ -null Borel set  $X \times 2^\omega$  still coded in  $\mathfrak{M}$ , so  $\langle a, b \rangle$  is not random.

If  $b$  is not random over  $\mathfrak{M}[a]$ , then  $b$  belongs to a  $\lambda$ -null set  $X$  coded in  $\mathfrak{M}[a]$ . By Borel reading of names,  $X$  has a Borel code of the form  $f(a)$ , where  $f : 2^\omega \rightarrow 2^\omega$  is a Borel map coded by some  $p \in 2^\omega \cap \mathfrak{M}$ , that is, a  $\Delta_1^1(p)$  map. This results in a  $\Delta_1^1(p)$  set  $P \subseteq 2^\omega \rightarrow 2^\omega$  such that  $\langle a, b \rangle \in P$  and the cross-section  $P_a = \{b' : \langle a, b' \rangle \in P\}$  (it contains  $b$ ) is  $\lambda$ -null. Therefore  $\langle a, b \rangle$  belongs to  $P' = \{\langle a', b' \rangle \in P : \lambda(P_{a'}) = 0\}$ , which is a  $\Pi_1^1(p)$  set (because being null is a  $\Pi_1^1$  property in this context by e.g. 2.2.3 in [7]) and a  $(\lambda \times \lambda)$ -null set by Fubini. Covering  $P'$  by a Borel null set coded in  $\mathfrak{M}$ , we conclude that  $\langle a, b \rangle$  is not random.

To prove that conversely (2) implies (1), suppose that (1) fails, that is, by (4),  $\langle a, b \rangle \in P$ , where  $P \subseteq 2^\omega \times 2^\omega$  is a  $(\lambda \times \lambda)$ -null  $\Delta_1^1(p)$  set,  $p \in 2^\omega \cap \mathfrak{M}$ . Consider the partition  $P = P' \cup P''$  of  $P$  into the  $\Pi_1^1(p)$  set  $P' = \{\langle a', b' \rangle \in P : \lambda(P_{a'}) = 0\}$  and the  $\Sigma_1^1(p)$  set  $P'' = P \setminus P'$ . Now if  $\langle a, b \rangle \in P'$ , then  $b$  belongs to the  $\lambda$ -null  $\Pi_1^1(p, a)$  set  $P'_a$ , and hence  $b$  is not random over  $\mathfrak{M}[a]$  (by covering  $P'_a$  by a Borel null set coded in  $\mathfrak{M}[a]$ ). If  $\langle a, b \rangle \in P''$ , then  $a$  belongs to the projection

$$\text{dom}(P'') = \{a' : \exists b' (\langle a', b' \rangle \in P)\} = \{a' : \lambda(P_{a'}) > 0\},$$

which is a  $\Sigma_1^1(p)$  set,  $\lambda$ -null by Fubini (as  $P$  is null), so  $a$  is not random. □

*Proof of Theorem 1, case B.* Similarly to case A, the contrary assumption leads to an ordinal  $1 \leq \xi < \omega_1$ , a code  $h \in \mathbf{FC}$  in  $\mathbf{V}$ , and a condition  $T_0 \in \mathbf{Rand}$  in  $\mathbf{V}$  which **Rand**-forces, over  $\mathbf{V}$ , that  $\vartheta_h(\mathfrak{a}) \in \mathbf{BC}_\xi$  and  $\{x \in 2^\omega : \varphi(x)\} = \mathbf{B}_{\vartheta_h(\mathfrak{a})}$ , where  $\mathfrak{a}$  is a canonical name for the random real and also forces that there is no Borel code  $c \in \mathbf{BC}_\xi$  in  $\mathbf{V}$ , satisfying  $\{x \in 2^\omega : \varphi(x)\} = \mathbf{B}_c$  in  $\mathbf{V}[\mathfrak{a}]$ .

**Argue in the universe  $\mathbf{V}$ .** There exists a closed non-null set  $D \subseteq [T_0]$  such that  $\vartheta_h(x) \in \mathbf{BC}_\xi$  for all  $x \in D$ . Consider the Borel set

$$P = \{\langle x, y \rangle : x \in D \wedge y \in \mathbf{B}_{\vartheta_h(x)}\} \subseteq 2^\omega \times \omega^\omega$$

and the  $\Pi_1^1$  equivalence relation  $x \mathbf{E} x'$  iff  $P_x = P_{x'}$ , on  $D$ , where  $P_x = \{y : P(x, y)\}$ . Then  $\mathbf{E}$  is  $\lambda$ -measurable, and so are all  $\mathbf{E}$ -equivalence classes. Thus there is a condition  $T_1 \in \mathbf{Rand}$  in  $\mathbf{V}$  which satisfies  $[T_1] \subseteq D$  and satisfies the requirements of one of the two cases:

*Case 1 (in  $\mathbf{V}$ ).* All  $\mathbf{E}$ -equivalence classes in  $[T_1]$  are  $\lambda$ -null sets. Then the  $\Pi_1^1$ -set  $W = \{\langle x, x' \rangle : x, x' \in [T_1] \wedge x \mathbf{E} x'\}$  is  $\lambda^2$ -null by Fubini. Therefore  $W$  is covered by a  $\mathbf{G}_\delta$  null set  $G \subseteq [T_1] \times [T_1]$ . Fix a transitive model  $\mathfrak{M} \in \mathbf{V}$  of a sufficiently large fragment of **ZFC** which contains  $h, T_1$ , and a code  $G$  and is an elementary submodel of the universe  $\mathbf{V}$  w.r.t. all analytic formulas.

**Lemma 5** (= Lemma 3.3 in [5]). *There are reals  $a, b \in [T_1]$ , separately random over  $\mathbf{V}$ , such that  $\mathbf{V}[a] = \mathbf{V}[b]$  and the pair  $\langle a, b \rangle$  is random over  $\mathfrak{M}$ .*

*Proof.* The set  $P = \{\langle x, x +_2 y \rangle : x, y \in [T_1]\}$  is non-null; hence, by Fubini, so is the projection  $Z = \{z \in 2^\omega : \lambda(P^z) > 0\}$ . Then follow the proof of Lemma 3, using Proposition 4 in treatment of the random pairs involved.  $\square$  (Lemma)  $\square$

The lemma leads to a contradiction similarly to Case 1 in Section 3.

*Case 2 (in  $\mathbf{V}$ ).* At least one of the  $\mathbf{E}$ -equivalence classes in  $[T_1]$  is not  $\lambda$ -null. Then, in  $\mathbf{V}$ , there is a condition  $T \subseteq T_1$  such that all points  $x \in [T]$  are  $\mathbf{E}$ -equivalent. In other words there exists a particular  $\Pi_\xi^0$  set  $A = \mathbf{B}_c$  with a code  $c \in \mathbf{BC}_\xi$  in  $\mathbf{V}$  such that  $\mathbf{B}_{\mathfrak{a}_h(x)} = \mathbf{B}_c$  for all  $x \in [T]$ . Then  $T$  forces over  $\mathbf{V}$  that  $\mathbf{B}_{\mathfrak{a}_h(\hat{\mathbf{a}})} = \mathbf{B}_c$ . But this contradicts the contrary assumption in the beginning of the proof, since  $T \subseteq T_0$  by construction.

$\square$  (Theorem 1, case B)  $\square$

## 5. SACKS REALS, CASE C OF THEOREM 1

The Sacks forcing  $\mathbf{Perf}$  consists of all perfect trees  $T \subseteq 2^{<\omega}$  (no endpoints and no isolated branches).

*Proof of Theorem 1, case C.* As above, the contrary assumption leads to an ordinal  $1 \leq \xi < \omega_1$ , a code  $h \in \mathbf{FC}$  in  $\mathbf{V}$ , and a condition  $T_0 \in \mathbf{Perf}$  in  $\mathbf{V}$  which  $\mathbf{Perf}$ -forces, over  $\mathbf{V}$ , that  $\mathfrak{a}_h(\hat{\mathbf{a}}) \in \mathbf{BC}_\xi$  and  $\{x \in 2^\omega : \varphi(x)\} = \mathbf{B}_{\mathfrak{a}_h(\hat{\mathbf{a}})}$ , where  $\hat{\mathbf{a}}$  is a canonical name for the Sacks-generic real, and also forces that there is no Borel code  $c \in \mathbf{BC}_\xi$  in  $\mathbf{V}$  satisfying  $\{x \in 2^\omega : \varphi(x)\} = \mathbf{B}_c$  in  $\mathbf{V}[\hat{\mathbf{a}}]$ .

**Argue in the universe  $\mathbf{V}$ .** There exists a condition  $T_1 \in \mathbf{Perf}$ ,  $T_1 \subseteq T_0$ , such that  $\mathfrak{a}_h(x) \in \mathbf{BC}_\xi$  for all  $x \in [T_1]$ . Consider the Borel set

$$P = \{\langle x, y \rangle : x \in [T_1] \wedge y \in \mathbf{B}_{\mathfrak{a}_h(x)}\} \subseteq 2^\omega \times \omega^\omega$$

and the  $\Pi_1^1$  equivalence relation  $x \mathbf{E} x'$  iff  $P_x = P_{x'}$ , on  $[T_1]$ . By the Silver dichotomy theorem,<sup>3</sup> there is a condition  $T \in \mathbf{Perf}$  in  $\mathbf{V}$  which satisfies  $T \subseteq T_1$  and satisfies the requirements of one of the two cases:

*Case 1 (in  $\mathbf{V}$ ).* The reals in  $[T]$  are pairwise  $\mathbf{E}$ -inequivalent. Then using any homeomorphism  $g : [T] \xrightarrow{\text{ontq}} [T]$  coded in  $\mathbf{V}$  and satisfying  $g(x) \neq x$  for all  $x$ , we easily get a pair of reals  $a \neq b$  in  $[T]$  Sacks generic over  $\mathbf{V}$  and satisfying  $\mathbf{V}[a] = \mathbf{V}[b]$ ; basically,  $b = g(a)$ . Then on the one hand,  $a \not\mathbf{E} b$  (since  $a \neq b$ ), thus  $P_a \neq P_b$ , that is,  $\mathbf{B}_{\mathfrak{a}_h(a)} \neq \mathbf{B}_{\mathfrak{a}_h(b)}$ , but on the other hand it is true in  $\mathbf{V}[a] = \mathbf{V}[b]$  that the sets  $\mathbf{B}_{\mathfrak{a}_h(a)}$  and  $\mathbf{B}_{\mathfrak{a}_h(b)}$  are equal to one and the same set  $\{x : \varphi(x)\}$ , which is a contradiction. Thus Case 1 is impossible.

*Case 2 ( $\mathbf{V}$ ).* The reals in  $[T]$  are pairwise  $\mathbf{E}$ -equivalent. Then, in  $\mathbf{V}$ , there is a particular  $\Pi_\xi^0$  set  $A = \mathbf{B}_c$  with a code  $c \in \mathbf{BC}_\xi$  in  $\mathbf{V}$  such that  $\mathbf{B}_{\mathfrak{a}_h(x)} = \mathbf{B}_c$  for all  $x \in [T]$ . Then  $T$  forces over  $\mathbf{V}$  that  $\mathbf{B}_{\mathfrak{a}_h(\hat{\mathbf{a}})} = \mathbf{B}_c$ . But this contradicts the contrary assumption in the beginning of the proof, since  $T \subseteq T_0$  by construction.

$\square$  (Theorem 1, case C)  $\square$

<sup>3</sup>Silver's theorem [11] claims that if  $\mathbf{E}$  is a  $\Pi_1^1$  equivalence relation on a Borel set  $X$ , then either there exist at most countably many  $\mathbf{E}$ -equivalence classes inside  $X$  or there is a perfect set  $Y \subseteq X$  of pairwise  $\mathbf{E}$ -inequivalent elements.

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